

Lax algebras in probability theory and topology

Ruben Van Belle

10th Workshop on Application of Categorical Probability Theory to Finance,
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Overview

- 1 Lax algebras
- 2 Lax algebras in topology ¹
- 3 Lax algebras in probability theory

¹*Monoidal topology*, D. Hofmann, G.J. Seal, W. Tholen

1. Lax algebras

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that induces the monad \mathbb{T} . This is the initial such adjunction.

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- 2 The algebras of the free group monad are *groups*.

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Together with 2-cells

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together with the appropriate coherence axioms.

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The category $\mathbf{2}\text{-Rel}$ is the usual 2-category of sets and relations.

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This fits in an entire theory of extensions of monads on \mathbf{Set} to lax monads on $V - \mathbf{Rel}$.

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Lax algebras vs. (strict) algebras

For a *large class*² of lax extensions of monads on **Set**, the inclusion functor

$$\mathbf{Set}^{\mathbb{T}} \rightarrow V\text{-Rel}_{\text{lax}}^{\mathbb{T}}$$

has a left adjoint.

²This includes all the lax extensions from section 1 and section 2. 

2. Topology

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By Zorn's lemma (AC), non-trivial ultrafilters exist.

Ultrafilters

Let X be a set. A subset $\mathcal{F} \subset \mathcal{P}X$ is called a **filter** on X if:

- $\mathcal{F} \neq \emptyset$,
- If $A, B \in \mathcal{F}$, then there is a $C \in \mathcal{F}$ such that $C \subseteq A$ and $C \subseteq B$,
- If $A \in \mathcal{F}$ and $A \subseteq B$, then $B \in \mathcal{F}$.

The collection of filters on X is partially ordered by inclusion. A maximal filter on X is called an *ultrafilter*.

Example: For $x \in X$, consider \mathcal{U}_x as the collection of all subsets of X that contain x . Then \mathcal{U}_x is a (trivial) ultrafilter.

By Zorn's lemma (AC), non-trivial ultrafilters exist.

Proposition

If \mathcal{U} is an ultrafilter on X and $A_1 \cup \dots \cup A_n \in \mathcal{U}$, then there is an n_0 such that $A_{n_0} \in \mathcal{U}$.

Ultrafilters in topology

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Example: Let $(x_n)_n$ be a sequence in X . Define \mathcal{F} as the filter generated by the set $(\{x_k \mid k \geq n\})_n$.

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We say that \mathcal{F} **converges** to x , if for every neighbourhood G of x , $G \in \mathcal{F}$.

Example: Let $(x_n)_n$ be a sequence in X . Define \mathcal{F} as the filter generated by the set $(\{x_k \mid k \geq n\})_n$. Then $(x_n)_n$ converges to x if and only if \mathcal{F} converges to x .

We write $\mathcal{F} \rightarrow x$

Proposition

If (X, \mathcal{T}) is Hausdorff, then every filter converges to at most one point.

Ultrafilters in topology

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If X is compact, then every ultrafilter on X converges to at least one point.

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Proof: Let \mathcal{U} be an ultrafilter on X . Suppose that \mathcal{U} does not converge, then for every $x \in X$ there is an open neighbourhood G_x such that $G_x \notin \mathcal{U}$.

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Proof: Let \mathcal{U} be an ultrafilter on X . Suppose that \mathcal{U} does not converge, then for every $x \in X$ there is an open neighbourhood G_x such that $G_x \notin \mathcal{U}$. Since X is compact, there is a finite collection x_1, \dots, x_n such that $X = \bigcup_{k=1}^n G_{x_k}$.

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The ultrafilter monad

For every set X , let $\beta(X)$ be the set of all ultrafilters on X .

³*Equational completion, model induced triples and pro-objects*, J.F.Kennison and D. Gildenhuys

The ultrafilter monad

For every set X , let $\beta(X)$ be the set of all ultrafilters on X .

For a map $f : X \rightarrow Y$ let βf be the map that sends an ultrafilter \mathcal{U} on X to the ultrafilter generated by $f(\mathcal{U})$.

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This defines an endofunctor $\beta : \mathbf{Set} \rightarrow \mathbf{Set}$.

- For a set X , define a map $\eta_X : X \rightarrow \beta X$ by the assignment

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$$\mu(\mathbf{U}) := \{A \mid \forall \mathcal{U} \in \mathbf{U} : A \in \mathcal{U}\}$$

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This monad is the codensity monad of the inclusion $\mathbf{Set}_f \rightarrow \mathbf{Set}$ of finite sets into all sets.³

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Algebras of the ultrafilter monad

Let X be a compact Hausdorff space.

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Theorem (Manes)

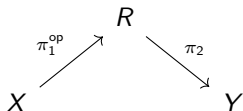
The category of algebras of the ultrafilter monad is equivalent to the category of compact Hausdorff spaces and continuous maps.

The ultrafilter monad on **Rel**

Let $R : X \rightsquigarrow Y$ be a relation of sets.

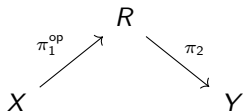
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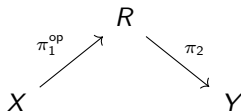
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This leads to the so called *Barr extensions* of the ultrafilter monad on **Set** to the ultrafilter monad on **Rel**.

Algebras of the ultrafilter monad on **Rel**

Let X be a topological space.

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This forms a lax algebras over β .

Theorem (Barr)

The category $\mathbf{Rel}_{\text{lax}}^\beta$ is equivalent to **Top**.

The Čech-Stone compactification

The inclusion $\mathbf{Set}^\beta \rightarrow \mathbf{Rel}_{\text{lax}}^\beta$ has a left adjoint.

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has a left adjoint. This is the *Čech-Stone compactification* functor.

3. Probability Theory

The distribution monads

For a set A , we can look at the set

$$DA := \left\{ (p_a)_{a \in A} \in [0, 1]^A \mid \text{finitely supported, } \sum_{a \in A} p_a = 1 \right\}.$$

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For a function $f : A \rightarrow B$, there is a function $Df : DA \rightarrow DB$ defined by the assignment:

$$(p_a)_{a \in A} \mapsto \left(\sum_{f(a)=b} p_a \right)_{b \in B}.$$

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This induces a functor $D : \mathbf{Set} \rightarrow \mathbf{Set}$.

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This gives D a monad structure. This monad is called the **distribution monad**.

Distribution monads

- Using the Barr extension construction, we can lift this monad to a monad on **Rel**, which we will also denote by D .
- If A is a metric space, topological space or measurable space, we can give DA extra structure, so that D becomes a monad on **Met**₁, **Top** or **Mble**.

Algebras of the distribution monads

A **convex set** is a set X together with a collection of operations $(c_\lambda : X \times X \rightarrow X)_{\lambda \in [0,1]}$ satisfying certain axioms.⁴

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The category **Set**^D is isomorphic to the category of convex sets and convex functions.

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Remark: For the distribution monads on **Met**₁, **Top** or **Mble**, it is *not* enough to have that the structure maps $(c_\lambda)_\lambda$ are in the respective category, because we also need continuity/measurability in the λ -variable.

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Proposition

The algebras of the distribution monads on **Met**₁ and **Top** are convex sets that *have* the cancellation property.

Lax algebras of the distribution monads

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Proposition

The category $\mathbf{Rel}_{\text{lax}}^D$ is isomorphic to the category of partial convex sets and partial convex functions.

Question: Does the functor $\mathbf{Set}^D \rightarrow \mathbf{Rel}_{\text{lax}}^D$ have a left adjoint?

Probability monads

Let X be a measurable space. Define GX as the set

{probability measures on X }

together with the σ -algebra generated by the evaluation maps

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This defines a map $\mu_X : GGX \rightarrow GX$, giving G a monad structure. This monad is called the **Giry monad**.⁵

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Probability monads

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Algebras of probability monads

Algebras are spaces where integration makes sense. These are the spaces where random variables should take their values.

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Proof(sketch):

$$\begin{array}{ccc} DX & \xrightarrow{\alpha} & X \\ \text{In} & \nearrow \tilde{\alpha} & \\ GX & & \end{array}$$

Every uniformly continuous function on a dense subset can be uniquely extended to the whole space.

Corollary

If X is an algebra for the distribution monad on \mathbf{Met}_1 , then \bar{X} is an algebra for the corresponding probability monad on \mathbf{Met}_1 .

Algebras of probability monads

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We would like to know when X itself is an algebra.

Theorem

Let X be an algebra for the distribution monad on \mathbf{Met}_1 . Suppose that

$$\forall x \in X \forall y \in \bar{X} \forall \lambda \in [0, 1] : \lambda x + (1 - \lambda)y \in \bar{X} \setminus X \Rightarrow \lambda = 0,$$

then X is an algebra for the corresponding probability monad.

Questions

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- 2 Can we construct lax extensions of monads on **Mble**, **Top**, **Met**₁, ...?
- 3 What would the lax algebras of these extended probability monads look like?
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- 2 Can we construct lax extensions of monads on **Mble**, **Top**, **Met**₁, ...?
- 3 What would the lax algebras of these extended probability monads look like?
- 4 Is there an adjunction between algebras and lax algebras? (Note that \mathbb{R} is not an algebra for the Giry monad)

- 1 Partial evaluations
- 2 Martingales in algebras



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