



Optimization of Truss Structures by Semidefinite Programming

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Outline

Introduction

Problem formulation

The optimization method

Techniques employed

Exploiting the algebraic structures Member adding Warm-start strategy

Numerical examples

Conclusions and future works

Structural optimization

• Consider the following design domain and loading conditions.



- The goal is to find the lightest structure that is able to carry the given set of loads.
- Several approaches of structural optimization.

Topology optimization (continuum*)

Topology optimization (truss)





*O. Sigmund. A 99 line topology optimization code written in Matlab. Structural and Multidisciplinary Optimization, 21:120–127, 2001.

Application



http://www.bbc.co.uk/programmes/p01rrnwc/p01rrbzl



https://en.wikipedia.org/wiki/Astoria%E2%80%93Megler_Bridge



https://en.wikipedia.org/wiki/London_King27s_Cross _railway_station



http://www.buildingtalk.com/wpcontent/uploads/arsenal-1.jpg

The underlying minimum weight problem

$$\begin{array}{ll} \underset{a,q_{\ell},u_{\ell}}{\text{minimize}} & I^{T} a \\ \text{subject to} & \sum_{i} q_{\ell,i} \gamma_{i} = f_{\ell}, \qquad \ell = 1, \cdots, n_{L} \\ & \frac{a_{i} E}{I_{i}} \gamma_{i}^{T} u_{\ell} = q_{\ell,i} \qquad \ell = 1, \cdots, n_{L}, i = 1, \cdots, m \\ & - a \sigma^{-} \leq q_{\ell} \leq \sigma^{+} a, \quad \ell = 1, \cdots, n_{L} \\ & a > 0 \end{array}$$

$$(1)$$

- n₁ number of load cases.
- ▶ $I \in \mathbb{R}^n$ is a vector of bar lengths,
- $a \in \mathbb{R}^n$ is a vector of bar cross-sectional areas,
- $f_{\ell} \in \mathbb{R}^m$ is a vector of applied load forces.
- ▶ $q_{\ell} \in \mathbb{R}^n$ are axial forces in members,
- $\sigma^- > 0$ and $\sigma^+ > 0$ are the the material's yield stresses in compression and tension,

E is Young's modulus. Bendsøe, M., Sigmund, O, Topology Optimization: Theory, Methods and Applications. Springer (2003) efinite Programming: Theory and Application, Edinburgh, 19 October 2018 A. Weldeysus, J. Gondzio



Stability constraints

• Consider the following three dimensional problem.







(a) Design domains, bc, and loads.

(b) Without stability considerations.

(c) With stability considerations.

Without stability considerations:

- The optimal design (a slender of six bars in compression) needs some kind of support or bracing from orthogonal directions.
- The optimal design for the bridge problem includes independent planar trusses. It lacks connectivity.

With stability considerations:

- The bar has bracing.
- The independent planar trusses in the bridge are connected

The minimum weight problem with global stability constraints minimize $I^T a$ a, q_{ℓ}, u_{ℓ} subject to $\sum_{i} q_{\ell,i} \gamma_i = f_{\ell}, \qquad \ell = 1, \cdots, n_L$ $\frac{a_i E}{l_i} \gamma_i^T u_\ell = q_{\ell,i} \qquad \ell = 1, \cdots, n_L, i = 1, \cdots, m$ (2) $-a\sigma^- < q_\ell < \sigma^+ a, \quad \ell = 1, \cdots, n_\ell$ $K(a) + \tau_{\ell} G(q_{\ell}) \succeq 0 \quad \ell = 1, \cdots, n_{\ell}$ a > 0

where the stiffness matrix K and geometry stiffness matrix G are given by

$$\mathcal{K}(a) = \sum_{j=1}^{m} a_j \mathcal{K}_j, \text{ with } \mathcal{K}_j = \frac{\mathcal{E}_j}{l_j} \gamma_j \gamma_j^{\mathsf{T}}, \text{ and } \mathcal{G}(q) = \sum_{j=1}^{m} q_j \mathcal{G}_j, \text{ with } \mathcal{G}_j = \frac{1}{l_j} (\delta_j \delta_j^{\mathsf{T}} + \eta_i \eta_j^{\mathsf{T}})$$

- The loading factor $\tau_{\ell} \geq 1$.
- $(\delta_j, \gamma_j, \eta_j)$ are mutually orthogonal. $(\eta = 0 \text{ for } 2D \text{ problems})$

M. Kocvara. On the modelling and solving of the truss design problem with global stability constraints. Structural and Multidisciplinary Optimization, 23:189–203, 2002. Semidefinite Programming: Theory and Application, Edinburgh, 19 October 2018 A. Weldeysus, J. Gondzio

The minimum weight problem with global stability constraints

$$\begin{array}{ll} \underset{a,q_{\ell},u_{\ell}}{\text{minimize}} & I^{T} a \\ \text{subject to} & \sum_{i} q_{\ell,i} \gamma_{i} = f_{\ell}, \qquad \ell = 1, \cdots, n_{L} \\ & \frac{a_{i}E}{l_{i}} \gamma_{i}^{T} u_{\ell} = q_{\ell,i} \qquad \ell = 1, \cdots, n_{L}, i = 1, \cdots, m \\ & - a\sigma^{-} \leq q_{\ell} \leq \sigma^{+} a, \quad \ell = 1, \cdots, n_{L} \\ & K(a) + \tau_{\ell} G(q_{\ell}) \succeq 0 \qquad \ell = 1, \cdots, n_{L} \\ & a \geq 0 \end{array}$$

$$(2)$$

The problem (2) is large-scale nonlinear non-convex semidefinite program.



Simplification

Ignore the kinematic compatibility constraints

$$\frac{a_i E}{l_i} \gamma_i^T u_\ell = q_{\ell,i}, \ \ell = 1, \cdots, n_L, \ i = 1, \cdots, m.$$

Hence, we solve the linear formulation

$$\begin{array}{ll} \underset{a,q_{\ell}}{\text{minimize}} & l^{T} a \\ \text{subject to} & \sum_{i} q_{\ell,i} \gamma_{i} = f_{\ell}, \qquad \ell = 1, \cdots, n_{L} \\ & -a\sigma^{-} \leq q_{\ell} \leq \sigma^{+} a, \quad \ell = 1, \cdots, n_{L} \\ & K(a) + \tau_{\ell} G(q_{\ell}) \succeq 0 \quad \ell = 1, \cdots, n_{L} \\ & a \geq 0. \end{array}$$

$$(3)$$

We then measure the violation due to ignoring the kinematic compatibility constraints by solving the least-squares problem

minimize
$$\max_{\ell} \frac{1}{||q_{\ell}^{*}||^{2}} \sum_{i} (\frac{a_{i}^{*}E}{l_{i}}\gamma_{i}^{T}u_{\ell} - q_{\ell,i}^{*})^{2},$$
 (4)

where a^* and q^*_{ℓ} are the solution of the relaxed problem (3). Semidefinite Programming: Theory and Application, Edinburgh, 19 October 2018 A. Weldeysus, J. Gondzio

Simplification

The (relaxation) SDP problem

$$\begin{array}{ll} \underset{a,q_{\ell}}{\text{minimize}} & I^{T} a \\ \text{subject to} & \sum_{i} q_{\ell,i} \gamma_{i} = f_{\ell}, \qquad \ell = 1, \cdots, n_{L} \\ & -a\sigma^{-} \leq q_{\ell} \leq \sigma^{+} a, \quad \ell = 1, \cdots, n_{L} \\ & K(a) + \tau_{\ell} G(q_{\ell}) \succeq 0 \quad \ell = 1, \cdots, n_{L} \\ & a \geq 0 \end{array}$$

$$(5)$$

- can be efficiently solved.
- provides lower bounds to the nonlinear and non-convex formulation.
- ► its solution has (usually) small violation in the kinematic compatibility constraints for realistic input and reasonable value of \(\tau_{\ell}\)

Primal-Dual Interior Point Method

Consider the following primal and dual semidefinite programs.

 $\begin{array}{cccc} & \underset{X}{\operatorname{minimize}} & \underset{C \bullet X}{\operatorname{minimize}} & \underset{C \bullet X}{\operatorname{minimize}} & \underset{y,S}{\operatorname{subject to}} & \underset{X \succeq 0}{\operatorname{minimize}} & \underset{y,S}{\operatorname{subject to}} & \underset{i=1}{\overset{m}{\operatorname{subject to}}} & \underset{i=$

Primal-Dual Interior Point Method

• Consider the following primal and dual semidefinite programs.

 $\max_{\substack{ \text{maximize} \\ maximize}} Dual_{b^{T}v}$ Primal minimize $C \bullet X$ subject to $\sum_{i=1}^{m} y_i A_i + S = C$ subject to $A_i \bullet X = b_i, i = 1, ..., m$ $X \succ 0$ $S \succ 0$ where $C, A_i \in \mathbb{S}^{n \times n}$, $b, y \in \mathbb{R}^m$, and $U \bullet V = \sum_i \sum_j U_{ij} V_{ij}$ for $U, V \in \mathbb{R}^{n \times n}$. • The first-order optimality conditions are (solved for $\mu_k \rightarrow 0$) AX = b $\mathcal{A}^* v + S = C$ (6) $X = \mu S^{-1}$.

where $\mathcal{A}: \mathbb{S}^n \to \mathbb{R}^m : \mathcal{A}X = (A_i \bullet X)_{i=1}^m$ and $\mathcal{A}^*: \mathbb{R}^m \to \mathbb{S}^n : \mathcal{A}^*y = \sum_{i=1}^m y_i A_i$

Primal-Dual Interior Point Method Consider the following primal and dual semidefinite programs. $\max_{\substack{\text{maximize}}} Dual b^T v$ Primal minimize $C \bullet X$ subject to $\sum_{i=1}^{n} y_i A_i + S = C$ subject to $A_i \bullet X = b_i, i = 1, ..., m$ $X \succ 0$ $S \succ 0$ where $C, A_i \in \mathbb{S}^{n \times n}$, $b, y \in \mathbb{R}^m$, and $U \bullet V = \sum_i \sum_i U_{ij} V_{ij}$ for $U, V \in \mathbb{R}^{n \times n}$. • The first-order optimality conditions are (solved for $\mu_k \rightarrow 0$) $\Delta X = h$ $\mathcal{A}^* \mathbf{v} + \mathbf{S} = \mathbf{C}$ (6)

 $X = \mu S^{-1}.$

where $\mathcal{A} : \mathbb{S}^n \to \mathbb{R}^m : \mathcal{A}X = (A_i \bullet X)_{i=1}^m$ and $\mathcal{A}^* : \mathbb{R}^m \to \mathbb{S}^n : \mathcal{A}^*y = \sum_{i=1}^m y_i A_i$ • Apply Newton method to solve (6).

$$\begin{bmatrix} 0 & \mathcal{A}^* & \mathcal{I} \\ \mathcal{A} & 0 & 0 \\ \mathcal{E} & 0 & \mathcal{F} \end{bmatrix} \begin{bmatrix} \Delta X \\ \Delta y \\ \Delta S \end{bmatrix} = \begin{bmatrix} \xi_d \\ \xi_p \\ \xi_c \end{bmatrix}.$$

where $\mathcal{E} = I \odot I$, $\mathcal{F} = X \odot S^{-1}$, and $P \odot Q : \mathbb{S}^n \to \mathbb{S}^n : (P \odot Q)U = \frac{1}{2}(PUQ^T) + QUP^T)$

Primal-Dual Interior Point Method

Consider the following primal and dual semidefinite programs.

 $\max_{\substack{ \text{maximize} \\ maximize}} Dual_{b^{T}v}$ Primal minimize $C \bullet X$ subject to $\sum_{i=1}^{m} y_i A_i + S = C$ subject to $A_i \bullet X = b_i, i = 1, ..., m$ $X \succ 0$ $S \succ 0$ where $C, A_i \in \mathbb{S}^{n \times n}$, $b, y \in \mathbb{R}^m$, and $U \bullet V = \sum_i \sum_i U_{ij} V_{ij}$ for $U, V \in \mathbb{R}^{n \times n}$. • The first-order optimality conditions are (solved for $\mu_k \rightarrow 0$) AX = b $\mathcal{A}^* v + S = C$ (6) $X = \mu S^{-1}$.

where $\mathcal{A} : \mathbb{S}^n \to \mathbb{R}^m : \mathcal{A}X = (A_i \bullet X)_{i=1}^m$ and $\mathcal{A}^* : \mathbb{R}^m \to \mathbb{S}^n : \mathcal{A}^*y = \sum_{i=1}^m y_i A_i$ • Usually solved for the reduced system (normal equations)

$$\mathcal{A}\mathcal{E}^{-1}\mathcal{F}\mathcal{A}^*\Delta y = -\xi_p + \mathcal{A}\mathcal{E}^{-1}(\xi_d - \mathcal{F}\xi_c).$$

Primal-Dual Interior Point Method

- The method obtains solution within modest number of iterations.
- Every iteration requires solving the linear system

$$\mathcal{A}\mathcal{E}^{-1}\mathcal{F}\mathcal{A}^*\Delta y = -\xi_{\rho} + \mathcal{A}\mathcal{E}^{-1}(\xi_d - \mathcal{F}\xi_c).$$

- Forming the system requires $O(mn^3 + m^2n^2)$ arithmetic operations (straightforward expressions) (bottle-neck)
- ▶ Large storage requirement (bottle-neck). $\mathcal{AE}^{-1}\mathcal{FA}^*$ is usually full matrix.

$$\begin{array}{ll} \underset{a,q_{\ell},u_{\ell}}{\text{minimize}} & l^{T}a \\ \text{subject to} & Bq_{\ell} = f_{\ell}, & \ell = 1, \cdots, n_{L} \\ & -a\sigma^{-} \leq q_{\ell} \leq \sigma^{+}a, \quad \ell = 1, \cdots, n_{L} \\ & K(a) + \tau_{\ell}G(q_{\ell}) \succeq 0 \quad \ell = 1, \cdots, n_{L} \\ & a \geq 0 \end{array}$$

$$(7)$$



Exploiting algebraic structures

The reduced system

$$\mathcal{A}\mathcal{E}^{-1}\mathcal{F}\mathcal{A}^*\Delta y = -\xi_{\rho} + \mathcal{A}\mathcal{E}^{-1}(\xi_d - \mathcal{F}\xi_c).$$

for the truss problem is

$$\begin{bmatrix} A_{11} & A_{12}^{T} & 0\\ A_{12} & A_{22} & \tilde{B}^{T}\\ 0 & \tilde{B} & 0 \end{bmatrix} \begin{bmatrix} \Delta a\\ \Delta q_{\ell}\\ \Delta \lambda_{\ell} \end{bmatrix} = \begin{bmatrix} \xi_{1}\\ \xi_{2}\\ \xi_{3} \end{bmatrix}, \text{ where}$$
(8)
$$(A_{11})_{ij} = -\sum_{\ell} X_{\ell} K_{i} S_{\ell}^{-1} \bullet K_{j} + (D_{11})_{ij}$$
$$(A_{12})_{ij} = -X_{\ell} K_{i} S_{\ell}^{-1} \bullet G_{j} + (D_{12})_{ij}, (A_{22})_{ij} = -X_{\ell} G_{i} S_{\ell}^{-1} \bullet G_{j} + (D_{22})_{ij}$$
$$K_{j} = \frac{E_{j}}{l_{j}} \gamma_{j} \gamma_{j}^{T}, \ G_{j} = \frac{1}{l_{j}} (\delta_{j} \delta_{j}^{T} + \eta_{i} \eta_{j}^{T}),$$

 D_{kl} diagonal matrices, and $U \bullet V = \sum_i \sum_j U_{ij} V_{ij}$ for $U, V \in \mathbb{R}^{n \times n}$.

Exploiting algebraic structures

$$\begin{bmatrix} A_{11} & A_{12}^{T} & 0\\ A_{12} & A_{22} & \tilde{B}^{T}\\ 0 & \tilde{B} & 0 \end{bmatrix} \begin{bmatrix} \Delta a\\ \Delta q_{\ell}\\ \Delta \lambda_{\ell} \end{bmatrix} = \begin{bmatrix} \xi_{1}\\ \xi_{2}\\ \xi_{3} \end{bmatrix}, \text{ where}$$
(8)
$$(A_{11})_{ij} = -\sum_{\ell} X_{\ell} K_{i} S_{\ell}^{-1} \bullet K_{j} + (D_{11})_{ij}$$
$$(A_{12})_{ij} = -X_{\ell} K_{i} S_{\ell}^{-1} \bullet G_{j} + (D_{12})_{ij}, (A_{22})_{ij} = -X_{\ell} G_{i} S_{\ell}^{-1} \bullet G_{j} + (D_{22})_{ij}$$
$$K_{j} = \frac{E_{j}}{I_{i}} \gamma_{j} \gamma_{j}^{T}, \ G_{j} = \frac{1}{I_{i}} (\delta_{j} \delta_{j}^{T} + \eta_{i} \eta_{j}^{T}),$$

• We exploit the low rank property (and sparsity) of the K_i 's and G_i 's.

$$X_{\ell} K_i S_{\ell}^{-1} \bullet K_j = \frac{E^2}{l_i l_j} \gamma_j^{\mathsf{T}} S_{\ell}^{-1} \gamma_i \gamma_i^{\mathsf{T}} X_{\ell} \gamma_j, \qquad (9)$$

▶ The matrix in (8) can be computed in $\mathcal{O}(n^2m)$ instead of $\mathcal{O}(nm^3 + n^2m^2)$ arithmetic operations.

A. Ben-Tal and A. Nemirovski. Robust truss topology design via semidefinite programming. SIAM Journal on Optimization, 7(4):991–1016, 1997.

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(a)



(b) Mem add iter.=1, 444 bars, $vol=0.05681m^3$



(d) Mem add iter.=3, 564 bars, $vol=0.05417m^3$



(f) Mem add iter.=5, 592 bars, vol= 0.05414m³

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(c) Mem add iter.=2, 518 bars, $vol= 0.05429m^3$



(e) Mem add iter.=4, 588 bars, $vol= 0.05414m^3$



(g) Mem add iter.=6, 600 bars, vol= $0.05414m^3_{A. Weldeysus, J. Gondzio}$

Member adding Primal (# inital bars) $\sum_{j\in\mathcal{K}_0}l_ja_j+\sum_{j\in\mathcal{K}_1}l_ja_j$

 $\underset{a,q}{\mathsf{minimize}}$

subject to

$$\sum_{j \in \mathcal{K}_0} \gamma_j q_{\ell,j} + \sum_{j \in \mathcal{K}_1} \gamma_j q_{\ell,j} = f_{\ell}, \qquad \forall \ell$$

$$-\sigma^{-}\mathbf{a}_{j} \leq q_{\ell,j} \leq \sigma^{+}\mathbf{a}_{j}, \qquad \qquad j \in \mathcal{K}_{0}, orall \ell$$

$$-\sigma^{-}\mathbf{a}_{j} \leq q_{\ell,j} \leq \sigma^{+}\mathbf{a}_{j}, \qquad \qquad j \in \mathcal{K}_{1}, orall \ell$$

$$\sum_{j \in \mathcal{K}_0} a_j K_j + \sum_{j \in \mathcal{K}_1} a_j K_j + \tau_{\ell} \sum_{j \in \mathcal{K}_0} q_{\ell,j} G_j + \tau_{\ell} \sum_{j \in \mathcal{K}_1} q_{\ell,j} G_j \succeq 0 \qquad \forall \ell$$

$$\mathbf{a}_{j} \geq \mathbf{0}, j \in \mathcal{K}_{0}, \mathbf{a}_{j} \geq \mathbf{0}, j \in \mathcal{K}_{1}, orall \ell$$

$$\sum_{\lambda,X}^{\bullet} \sum_{\ell}^{\mathsf{Dual}} f_{\ell}^{\mathsf{T}} \lambda_{\ell}$$
s.t.
$$-\frac{I}{\sigma^{-}} (I_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}) \leq \sum_{\ell} (\gamma_{j}^{\mathsf{T}} \lambda_{\ell} + \tau_{\ell} G_{j} \bullet X_{\ell}) \leq \frac{1}{\sigma^{+}} (I_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}), j \in \mathcal{K}_{0}, \forall \ell$$

$$-\frac{I}{\sigma^{-}} (I_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}) \leq \sum_{\ell} (\gamma_{j}^{\mathsf{T}} \lambda_{\ell} + \tau_{\ell} G_{j} \bullet X_{\ell}) \leq \frac{1}{\sigma^{+}} (I_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}), j \in \mathcal{K}_{1}, \forall \ell$$

 $X_{\ell} \succeq 0, \forall \ell.$

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(# all bars)



- Set a = 0, q = 0.
- Primal

minimize $\sum_{j \in \mathcal{K}_{n}} l_{j} a_{j} + \sum_{i \in \mathcal{K}_{n}} l_{j} a_{j}^{*}$ subject to $\sum_{j \in \mathcal{K}_0} \gamma_j q_{\ell,j} + \sum_{j \in \mathcal{K}_1} \gamma_j q_{\ell,j}^{*,0} = f_{\ell},$ ∀ℓ $-\sigma^{-}a_{i} \leq q_{\ell,i} \leq \sigma^{+}a_{i},$ $i \in \mathcal{K}_0, \forall \ell$ $\sum_{j \in \mathcal{K}_0} a_j \mathcal{K}_j + \sum_{j \in \mathcal{K}_1} a_j \mathcal{K}_j^{\vee} + \tau_\ell \sum_{j \in \mathcal{K}_0} q_{\ell,j} \mathcal{G}_j + \tau_\ell \sum_{j \in \mathcal{K}_1} q_{\ell,j} \mathcal{G}_j^{\vee} \succeq 0$ ∀ℓ $a_j \geq 0, j \in \mathcal{K}_0, a_j \geq 0, j$







Solve the primal restricted problem (RMP)

$$\begin{array}{ll} \underset{a,q}{\text{minimize}} & \sum_{j \in \mathcal{K}_0} l_j a_j \\ \text{subject to} & \sum_{j \in \mathcal{K}_0} \gamma_j q_{\ell,j} = f_\ell, & \forall \ell \\ & -\sigma^- a_j \leq q_{\ell,j} \leq \sigma^+ a_j, & j \in \mathcal{K}_0, \forall \ell \\ & \sum_{j \in \mathcal{K}_0} a_j K_j + \tau_\ell \sum_{j \in \mathcal{K}_0} q_j G_j \succeq 0, & \forall \ell \\ & a_j \geq 0, & j \in \mathcal{K}_0 \end{array}$$

and the dual restricted problem (D-RMP)

$$\max_{\lambda, X} \sum_{\ell} f_{\ell}^{T} \lambda_{\ell}$$
s.t. $-\frac{1}{\sigma^{-}} (l_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}) \leq \sum_{\ell} (\gamma_{j}^{T} \lambda_{\ell} + \tau_{\ell} G_{j} \bullet X_{\ell}) \leq \frac{1}{\sigma^{+}} (l_{j} - \sum_{\ell} K_{j} \bullet X_{\ell}), j \in \mathcal{K}_{\ell}$
Semidefinite Programming: Theory and Application, Edinburgh, ¹⁹ October 2018 A. Weldeysus, J. Gondzio 22

• Generate the columns(matrices) as below.

$$\mathcal{K} = \left\{ j \in \{1, \cdots, m\} \setminus \mathcal{K}_0 | \sum_{\ell} (\gamma_j^T \lambda_{\ell}^* + \tau_{\ell} \mathcal{G}_j \bullet X_{\ell}^*) < -\frac{l}{\sigma^-} (l_j - \sum_{\ell} \mathcal{K}_j \bullet X_{\ell}^*) \text{ or } \right.$$
$$\left. \sum_{\ell} (\gamma_j^T \lambda_{\ell}^* + \tau_{\ell} \mathcal{G}_j \bullet X_{\ell}^*) > \frac{l}{\sigma^+} (l_j - \sum_{\ell} \mathcal{K}_j \bullet X_{\ell}^*) \right\}$$
(10)

where λ_{ℓ}^* and X_{ℓ}^* are solution of the D-RMPs.

- Filter, add, and the next problem instance.
- The sparsity of K_i 's and G_i is exploited to generate the set \mathcal{K}



(a) Mem add iter.=1, 444 bars, vol= $0.05681m^3$



(b) Mem add iter.=2, 518 bars, $vol=0.05429m^3$

Figure

• Generate the columns(matrices) as below.

$$\mathcal{K} = \left\{ j \in \{1, \cdots, m\} \setminus \mathcal{K}_0 | \sum_{\ell} (\gamma_j^T \lambda_{\ell}^* + \tau_{\ell} \mathbf{G}_j \bullet \mathbf{X}_{\ell}^*) < -\frac{l}{\sigma^-} (l_j - \sum_{\ell} \mathbf{K}_j \bullet \mathbf{X}_{\ell}^*) \text{ or } \right.$$
$$\left. \sum_{\ell} (\gamma_j^T \lambda_{\ell}^* + \tau_{\ell} \mathbf{G}_j \bullet \mathbf{X}_{\ell}^*) > \frac{l}{\sigma^+} (l_j - \sum_{\ell} \mathbf{K}_j \bullet \mathbf{X}_{\ell}^*) \right\}$$
(10)

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(a) Mem add iter.=1, 444 bars, vol= $0.05681m^3$



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Figure

Member adding the primal restricted problem (RMP) minimize $\sum_{a,q} l_j a_j + \sum_{l \neq a_j} l_j a_j$ $i \in K_0$ $i \in K$ subject to $\sum \gamma_j q_{\ell,j} + \sum \gamma_j q_{\ell,j} = f_{\ell}$, ∀ℓ $i \in \mathcal{K}_0$ $i \in \mathcal{K}$ $-\sigma^{-}a_{i} < q_{\ell,i} < \sigma^{+}a_{i}$ $i \in \mathcal{K}_0, \forall \ell$ $-\sigma^{-}a_{i} \leq q_{\ell,i} \leq \sigma^{+}a_{i},$ $i \in \mathcal{K}, \forall \ell$ $\sum_{j \in \mathcal{K}_0} a_j \mathcal{K}_j + \sum_{j \in \mathcal{K}} a_j \mathcal{K}_j + \tau_{\ell} \sum_{j \in \mathcal{K}_0} q_{\ell,j} \mathcal{G}_j + \tau_{\ell} \sum_{j \in \mathcal{K}} q_{\ell,j} \mathcal{G}_j \succeq 0$ ∀ℓ $a_i > 0, j \in \mathcal{K}_0, a_i > 0, j \in \mathcal{K}, \forall \ell$. the dual restricted problem (D-RMP) $\max_{\lambda,X} \quad \sum f_{\ell}^{T} \lambda_{\ell}$ s.t. $-\frac{l}{\sigma^-}(l_j - \sum_{i} K_j \bullet X_\ell) \leq \sum_{i} (\gamma_j^\mathsf{T} \lambda_\ell + \tau_\ell G_j \bullet X_\ell) \leq \frac{1}{\sigma^+}(l_j - \sum_{i} K_j \bullet X_\ell), j \in \mathcal{K}_0, \forall \ell$ $-\frac{l}{\sigma^+}(l_j-\sum_i K_j \bullet X_\ell) \leq \sum_i (\gamma_j^T \lambda_\ell + \tau_\ell \mathsf{G}_j \bullet X_\ell) \leq \frac{1}{\sigma^+}(l_j-\sum_i K_j \bullet X_\ell), j \in \mathcal{K}, \forall \ell$ $X_{\ell} \succ 0. \ \forall \ell.$

Warm-start strategy



(a)







(d) Mem add iter.=3, 564 bars, $vol=0.05417m^3$



(f) Mem add iter.=5, 592 bars, vol= 0.05414m³

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(c) Mem add iter.=2, 518 bars, $vol= 0.05429m^3$



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(g) Mem add iter.=6, 600 bars, vol= $0.05414m_{A. Weldeysus, J. Gondzio}^{3}$

Warm-start strategy





We extend the warm-start strategy

$$\begin{aligned} (\mathsf{a}, \mathsf{q}_{\ell}, \mathsf{S}_{\ell}, \mathsf{s}_{\ell}^+, \mathsf{s}_{\ell}^-) &\to (\mathsf{a}, \bar{\mathsf{a}}, \mathsf{q}_{\ell}, \bar{\mathsf{q}}_{\ell}, \mathsf{S}_{\ell}, \mathsf{s}_{\ell}^+, \bar{\mathsf{s}}_{\ell}^+, \mathsf{s}_{\ell}^-, \bar{\mathsf{s}}_{\ell}^-) \\ (\lambda_{\ell}, X_{\ell}, x_{\ell}^+, x_{\ell}^-) &\to (\lambda_{\ell}, X_{\ell}, x_{\ell}^+, \bar{x}_{\ell}^+, x_{\ell}^-, \bar{x}_{\ell}^-) \end{aligned}$$

The variables with the super-bar are vectors in \mathbb{R}^k , k = |K|

- Computing a warm-start point
 - \blacktriangleright Old variables \leftarrow solution of the preceding save problem with loose tolerance, say $\epsilon_{opt}=0.1$
 - New variables (those with super-bar)

$$\begin{split} & \bar{x}_{\ell,j}^{+} = \max\{\bar{\gamma}_{j}^{T} \lambda_{\ell} + \tau_{\ell} \bar{G}_{i} \bullet X_{\ell}, \mu_{0}^{\frac{1}{2}}\}, \, \forall j \in K, \\ & \bar{x}_{\ell,j}^{-} = \max\{-\bar{\gamma}_{j}^{T} \lambda_{\ell} - \tau_{\ell} \bar{G}_{i} \bullet X_{\ell}, \mu_{0}^{\frac{1}{2}}\}, \, \forall j \in K, \\ & \bullet (\bar{x}_{a})_{j} = \max\{|\bar{l}_{j} - \sigma^{+} \sum_{\ell} (\bar{x}^{+}_{\ell})_{j} - \sigma^{-} \sum_{\ell} (\bar{x}^{-}_{\ell})_{j} - \bar{K}_{i} \bullet X_{\ell}|, \mu_{0}^{\frac{1}{2}}\}, \, \forall j \in K, \\ & \bullet \bar{q}_{\ell}^{+} = 0 \, \forall \ell \in \{1, ..., n_{L}\} \\ & \bar{a}_{j} = \mu(\bar{x}_{a}^{-1})_{j}, \, \forall j \in K, \\ & \bullet \bar{s}_{\ell}^{+} = \sigma^{+}\bar{a}, \, \forall \ell \in \{1, ..., n_{L}\} \\ & \bullet \bar{s}_{\ell}^{-} = \sigma^{-}\bar{a}, \, \forall \ell \in \{1, ..., n_{L}\} \end{split}$$

Violation estimations

• Primal infeasibility $(\xi_{P_{1,\ell}}, \xi_{P_{2,\ell}}, \xi_{P_{3,\ell}}, \xi_{P_{4,\ell}})$

$$\begin{aligned} ||\xi_{p_{1,\ell}}||_{\infty} &= ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i} - \sum_{i} \bar{q}_{\ell,i}\bar{\gamma}_{i}||_{\infty} = ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i}||_{\infty} = ||\xi_{p_{1,\ell}}^{0}||_{\infty}, \\ ||\xi_{p_{2,\ell}}||_{\infty} &= ||\sigma^{+}\bar{a} - \bar{q}_{\ell} - \bar{s}_{\ell}^{+}||_{\infty} = 0, \\ ||\xi_{p_{3,\ell}}||_{\infty} &= ||\sigma^{-}\bar{a} + \bar{q}_{\ell} - \bar{s}_{\ell}^{-}||_{\infty} = 0, \\ ||\xi_{p_{4,\ell}}||_{\infty} &= || - K(a) - \tau_{\ell}G(q_{\ell}) + S_{\ell} - \bar{K}(\bar{a}) - \tau_{\ell}\bar{G}(\bar{q}_{\ell})||_{\infty} \leq ||\xi_{p_{4,\ell}}^{0}||_{\infty} + \mu_{0}^{\frac{1}{2}} \sum \frac{E_{i}}{\bar{I}_{i}} \end{aligned}$$
(11)

Violation estimations

▶ Primal infeasibility $(\xi_{P_{1,\ell}}, \xi_{P_{2,\ell}}, \xi_{P_{3,\ell}}, \xi_{P_{4,\ell}})$

$$\begin{aligned} ||\xi_{P_{1,\ell}}||_{\infty} &= ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i} - \sum_{i} \bar{q}_{\ell,i}\bar{\gamma}_{i}||_{\infty} = ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i}||_{\infty} = ||\xi_{P_{1,\ell}}^{0}||_{\infty}, \\ ||\xi_{P_{2,\ell}}||_{\infty} &= ||\sigma^{+}\bar{a} - \bar{q}_{\ell} - \bar{s}_{\ell}^{+}||_{\infty} = 0, \\ ||\xi_{P_{3,\ell}}||_{\infty} &= ||\sigma^{-}\bar{a} + \bar{q}_{\ell} - \bar{s}_{\ell}^{-}||_{\infty} = 0, \\ ||\xi_{P_{4,\ell}}||_{\infty} &= || - K(a) - \tau_{\ell}G(q_{\ell}) + S_{\ell} - \bar{K}(\bar{a}) - \tau_{\ell}\bar{G}(\bar{q}_{\ell})||_{\infty} \leq ||\xi_{P_{4,\ell}}^{0}||_{\infty} + \mu_{0}^{\frac{1}{2}} \sum_{i} \frac{E_{i}}{\bar{I}_{i}} \end{aligned}$$
(11)

Dual infeasibility $(\xi_{d_1}, \xi_{d_2,\ell})$

$$\begin{split} ||\xi_{d_{1}}||_{\infty} &= ||\sum_{\ell} (\sigma^{+} x_{\ell}^{+} + \sigma^{-} x_{\ell}^{-} + \mathcal{K} X_{\ell}) + x_{s} - I||_{\infty} \\ &\leq ||2(\bar{I} + \sum_{\ell} (\sigma_{max}(\varepsilon_{\ell}^{-} + \varepsilon_{\ell}^{+}) - \bar{\mathcal{K}} X_{\ell})) + (4n_{L} + 1)\mu_{0}^{\frac{1}{2}} e||_{\infty} \\ &||\bar{\xi}_{d_{2},\ell}||_{\infty} &= ||\bar{B}^{T} \lambda_{\ell} - \bar{x}_{\ell}^{+} + \bar{x}_{\ell}^{-} + \tau_{\ell} \bar{\mathcal{G}} X_{\ell}|| \leq \mu_{0}^{\frac{1}{2}}. \end{split}$$

Violation estimations

• Primal infeasibility $(\xi_{P_{1,\ell}}, \xi_{P_{2,\ell}}, \xi_{P_{3,\ell}}, \xi_{P_{4,\ell}})$

$$\begin{aligned} ||\xi_{P_{1,\ell}}||_{\infty} &= ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i} - \sum_{i} \bar{q}_{\ell,i}\bar{\gamma}_{i}||_{\infty} = ||f_{\ell} - \sum_{i} q_{\ell,i}\gamma_{i}||_{\infty} = ||\xi_{P_{1,\ell}}^{0}||_{\infty}, \\ ||\xi_{P_{2,\ell}}||_{\infty} &= ||\sigma^{+}\bar{a} - \bar{q}_{\ell} - \bar{s}_{\ell}^{+}||_{\infty} = 0, \\ ||\xi_{P_{3,\ell}}||_{\infty} &= ||\sigma^{-}\bar{a} + \bar{q}_{\ell} - \bar{s}_{\ell}^{-}||_{\infty} = 0, \\ ||\xi_{P_{4,\ell}}||_{\infty} &= || - K(a) - \tau_{\ell}G(q_{\ell}) + S_{\ell} - \bar{K}(\bar{a}) - \tau_{\ell}\bar{G}(\bar{q}_{\ell})||_{\infty} \leq ||\xi_{P_{4,\ell}}^{0}||_{\infty} + \mu_{0}^{\frac{1}{2}} \sum \frac{E_{i}}{\bar{l}_{i}} \end{aligned}$$
(11)

1

30

Dual infeasibility $(\xi_{d_1}, \xi_{d_{2,\ell}})$

$$\begin{aligned} ||\xi_{d_1}||_{\infty} &\leq ||2(\bar{l} + \sum_{\ell} (\sigma_{\max}(\varepsilon_{\ell}^- + \varepsilon_{\ell}^+) - \bar{\mathcal{K}}X_{\ell})) + (4n_L + 1)\mu_0^{\frac{1}{2}} e||_{\infty} \\ ||\bar{\xi}_{d_2,\ell}||_{\infty} &= ||\bar{B}^T \lambda_{\ell} - \bar{x}_{\ell}^+ + \bar{x}_{\ell}^- + \tau_{\ell}\bar{\mathcal{G}}X_{\ell}|| \leq \mu_0^{\frac{1}{2}}. \end{aligned}$$

 $\begin{array}{||c||} \bullet \quad \textbf{Centrality} \; (\bar{a}, \bar{s}_a), \; (\bar{X}_\ell, \bar{S}_\ell), (\bar{x}_\ell^-, \bar{s}_\ell^-), (\bar{x}_\ell^+, \bar{s}_\ell^+) \\ (\bar{a}, \bar{s}_a), \; (\bar{X}_\ell, \bar{S}_\ell) \; \text{are} \; \mu_0 \; \text{centered.} \end{array}$

$$\begin{aligned} \frac{\sigma^+}{\sigma_{\max}n_L\mu_0^{-\frac{1}{2}}(\max_{\ell}(\varepsilon_{\ell_j}^-+\varepsilon_{\ell_j}^+)+\bar{K}_i\bullet X_{\ell})+2n_L}\mu_0 &\leq (\bar{x}_{\ell}^+)_j(\bar{s}_{\ell}^+)_j \leq \mu_0\sigma^++\mu_0^{\frac{1}{2}}\sigma^+(\varepsilon_{\ell_j}^-+\varepsilon_{\ell_j}^+), \;\forall j, \forall \ell. \\ \frac{\sigma^-}{\sigma_{\max}n_L\mu_0^{-\frac{1}{2}}(\max_{\ell}(\varepsilon_{\ell_j}^-+\varepsilon_{\ell_j}^+)+\bar{K}_i\bullet X_{\ell})+2n_L}\mu_0 &\leq (\bar{x}_{\ell}^-)_j(\bar{s}_{\ell}^-)_j \leq \mu_0\sigma^-+\mu_0^{\frac{1}{2}}\sigma^-(\varepsilon_{\ell_j}^-+\varepsilon_{\ell_j}^+), \;\forall j, \forall \ell. \end{aligned}$$
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minimize	Nonlinear I ^T a			Relaxation	
subject to	$\sum q_{\ell,i} \gamma_i = f_\ell,$	$\forall \ell$	$\underset{a,q_{\ell},u_{\ell}}{\text{minimize}}$	l ^T a	
	i a _i E _T	\/ <i>0</i> \/'	subject to	$\sum_{i} q_{\ell,i} \gamma_i = f_\ell,$	$\forall \ell$
	$\overline{I_i} \gamma_i^{*} u_{\ell} = q_{\ell,i}$	$\forall \ell, \forall I$		$-a\sigma^{-}\leq q_{\ell}\leq \sigma^{+}a,$	$\forall \ell$
	$-a\sigma \leq q_{\ell} \leq \sigma \cdot a,$ $K(a) + \tau_{\ell} G(q_{\ell}) \succ 0$	$\forall \ell$ $\forall \ell$		$K(a) + au_\ell G(q_\ell) \succeq 0$	$\forall \ell$
	$a \ge 0$			$a \ge 0$	

Least-squares (LSQ) problem

minimize
$$\max_{\ell} \frac{1}{||q_{\ell}^{*}||^{2}} \sum_{i} (\frac{a_{i}^{*}E}{I_{i}}\gamma_{i}^{T}u_{\ell} - q_{\ell,i}^{*})^{2},$$
 (12)







au	0	1	10
Volume (nonlinear SDP)	0.062	0.06222	0.06464
Volume (relaxed SDP)	-	0.06217	0.06433
Violation (LSQ problem)	-	4.96e-06	5.32e-4

M. Kocvara. On the modelling and solving of the truss design problem with global stability constraints. Structural and Multidisciplinary Optimization, 23:189–203, 2002.

Semidefinite Programming: Theory and Application, Edinburgh, 19 October 2018

A. Weldeysus, J. Gondzio



Figure: $\tau_{\ell} = 20, 30, 40, ..., 90$. (a)-(h) By solving the relaxation linear SDP. (i)-(p) By solving the nonlinear SDP.

τ	20	30	40	50	60	70	80	90
Volume (nonlinear SDP)	0.0677	0.0717	0.0772	0.0846	0.0933	0.1031	0.1139	0.1251
Volume (relaxed SDP)	0.0670	0.0703	0.0749	0.0805	0.0871	0.0947	0.1028	0.1117
Violation (LSQ problem)	0.0024	0.0066	0.0164	0.0306	0.0368	0.0459	0.0591	0.0702

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A. Weldeysus, J. Gondzio



τ	0	1	10
Volume (nonlinear SDP)	0.0300	0.0302	0.0320
Volume (relaxed SDP)	-	0.0301	0.0310
Violation (LSQ problem)	-	3.7e-6	5.1e-5



Figure: Optimal design with stability constraints for $\tau_{\ell} = 20, 30, 40$. (a)-(c) By solving the linear SDP relaxation. (d)-(f) By solving the nonlinear SDP.

τ	20	30	40
Volume (nonlinear SDP)	0.0370	0.0507	0.0663
Volume (relaxed SDP)	0.0358	0.0499	0.0642
Violation (LSQ problem)	0.0151	0.0510	0.5889

Example: Validating the member adding

	All at once	With member adding
Volume (m ³)	0.05414	0.05414
Final number of bars	3240	600
Mem. add. iter	1	6
Total CPU (Sec)	145	28



The violation of the compatibility constraints by stable design is equal to 5.8336e - 06.

Example: Large-scale truss problems (SDP)

3638

- ▶ #bars =90,100 (*n* = 1,263, *m* = 180,200 in standard SDP notation).
- The full-scale SDP requires at least 240GB memory.

Total CPU



The violation of the compatibility constraints by stable design is equal to 52354e - 06. Semidefinite Programming: Theory and Application, Edinburgh, 19 October 2018 A. Weldeysus, J. Gondzio 39

2654

Example: Stadium roof (multiple-load cases)

36856 members, 3-loads case, 2487 bars and 6 mem add iter needed. CPU=2238Sec.



Figure: f_1 (red), f_2 (blue), and f_3 (green). A = (0, 0, 2.3), B = (5, 0, 0), C = (15, 0, 0), D = (20, 0, 0), E = (40, 0, 2.8), F = (15, 0, 2.4.2). The roof is 80m the y-direction.

The violation of the compatibility constraints by unstable design is equal to 0.0011. The violation of the compatibility constraints by stable design is equal to 0.3190.

Conclusions and future works

Conclusions

- Extended the member adding procedure to SDP.
- Developed and implemented a specialized primal-dual interior point method The method and its implementation:
 - exploits the structure of the problem.
 - uses warm-start strategy.

Future work

- Comparison to other SDP solvers.
- Look into the possibilities of using iterative methods for solving the linear systems.

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Thank you for your attention!