

Functional equations

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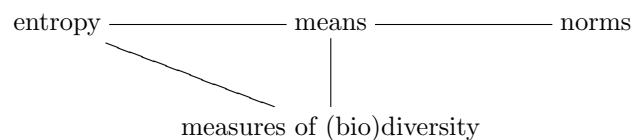
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Preamble

Hello.

Admin: email addresses; sections in outline \neq lectures; pace.

Overall plan: interested in unique characterizations of ...



There are many ways to measure diversity: long controversy.

Ideal: be able to say ‘If you want your diversity measure to have properties X, Y and Z, then it must be one of the following measures.’

Similar results have been proved for entropy, means and norms.

This is a tiny part of the field of functional equations!

One ulterior motive: for me to learn something about FEs. I’m not an expert, and towards end, this will get to edge of research (i.e. I’ll be making it up as I go along).

Tools:

- native wit
- elementary real analysis
- (new!) some probabilistic methods.

One ref: Aczél and Daróczy, *On Measures of Information and Their Characterizations*. (Comments.) Other refs: will give as we go along.

1 Warm-up

Week I (7 Feb)

Which functions f satisfy $f(x + y) = f(x) + f(y)$? Which functions of two variables can be separated as a product of functions of one variable?

This section is an intro to basic techniques. We may or may not need the actual results we prove.

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The Cauchy functional equation

The Cauchy FE on a function $f: \mathbb{R} \rightarrow \mathbb{R}$ is

$$\forall x, y \in \mathbb{R}, \quad f(x + y) = f(x) + f(y). \quad (1)$$

There are some obvious solutions. Are they the only ones? Weak result first, to illustrate technique.

Proposition 1.1 *Let $f: \mathbb{R} \rightarrow \mathbb{R}$ be a differentiable function. TFAE (the following are equivalent):*

- i. f satisfies (1)
- ii. there exists $c \in \mathbb{R}$ such that

$$\forall x \in \mathbb{R}, \quad f(x) = cx.$$

If these conditions hold then $c = f(1)$.

Proof (ii) \Rightarrow (i) and last part: obvious.

Now assume (i). Differentiate both sides of (1) with respect to x :

$$\forall x, y \in \mathbb{R}, \quad f'(x + y) = f'(x).$$

Take $x = 0$: then $f'(y) = f'(0)$ for all $y \in \mathbb{R}$. So f' is constant, so there exist $c, d \in \mathbb{R}$ such that

$$\forall x \in \mathbb{R}, \quad f(x) = cx + d.$$

Substituting back into (1) gives $d = 0$, proving (ii). □

‘Differentiable’ is a much stronger condition than necessary!

Theorem 1.2 *As for Proposition 1.1, but with ‘continuous’ in place of ‘differentiable’.*

Proof Let f be a continuous function satisfying (1).

- $f(0 + 0) = f(0) + f(0)$, so $f(0) = 0$.
- $f(x) + f(-x) = f(x + (-x)) = f(0) = 0$, so $f(-x) = -f(x)$. Cf. group homomorphisms.
- Next, $f(nx) = nf(x)$ for all $x \in \mathbb{R}$ and $n \in \mathbb{Z}$. For $n > 0$, true by induction. For $n = 0$, says $f(0) = 0$. For $n < 0$, have $-n > 0$ and so $f(nx) = -f(-nx) = -(-nf(x)) = nf(x)$.
- In particular, $f(n) = nf(1)$ for all $n \in \mathbb{Z}$.
- For $m, n \in \mathbb{Z}$ with $n \neq 0$, we have

$$f(n \cdot m/n) = f(m) = mf(1)$$

but also

$$f(n \cdot m/n) = nf(m/n),$$

so $f(m/n) = (m/n)f(1)$. Hence $f(x) = f(1)x$ for all $x \in \mathbb{Q}$.

- Now f and $x \mapsto f(1)x$ are continuous functions on \mathbb{R} agreeing on \mathbb{Q} , hence are equal. \square

Remarks 1.3 i. ‘Continuous’ can be relaxed further still. It was pointed out in class that continuity at 0 is enough. ‘Measurable’ is also enough (Fréchet, ‘Pri la funkcio $f(x+y) = f(x) + f(y)$ ’, 1913). Even weaker: ‘bounded on some set of positive measure’. But never mind! For this course, I’ll be content to assume continuity.

- ii. To get a ‘weird’ solution of Cauchy FE (i.e. not of the form $x \mapsto cx$), need existence of a non-measurable function. So, need some form of choice. So, can’t really construct one.
- iii. Assuming choice, weird solutions exist. Choose basis B for the vector space \mathbb{R} over \mathbb{Q} . Pick $b_0 \neq b_1$ in B and a function $\phi: B \rightarrow \mathbb{R}$ such that $\phi(b_0) = 0$ and $\phi(b_1) = 1$. Extend to \mathbb{Q} -linear map $f: \mathbb{R} \rightarrow \mathbb{R}$. Then $f(b_0) = 0$ with $b_0 \neq 0$, but $f \neq 0$ since $f(b_1) = 1$. So f cannot be of the form $x \mapsto cx$. But f satisfies the Cauchy functional equation, by linearity.

Variants (got by using the group isomorphism $(\mathbb{R}, +) \cong ((0, \infty), 1)$ defined by \exp and \log):

Corollary 1.4 i. Let $f: \mathbb{R} \rightarrow (0, \infty)$ be a continuous function. TFAE:

- $f(x+y) = f(x)f(y)$ for all x, y
- there exists $c \in \mathbb{R}$ such that $f(x) = e^{cx}$ for all x .

ii. Let $f: (0, \infty) \rightarrow \mathbb{R}$ be a continuous function. TFAE:

- $f(xy) = f(x) + f(y)$ for all x, y
- there exists $c \in \mathbb{R}$ such that $f(x) = c \log x$ for all x .

iii. Let $f: (0, \infty) \rightarrow (0, \infty)$ be a continuous function. TFAE:

- $f(xy) = f(x)f(y)$ for all x, y
- there exists $c \in \mathbb{R}$ such that $f(x) = x^c$ for all x .

Proof For (i), define $g: \mathbb{R} \rightarrow \mathbb{R}$ by $g(x) = \log f(x)$. Then g is continuous and satisfies Cauchy FE, so $g(x) = cx$ for some constant c , and then $f(x) = e^{cx}$.

(ii) and (iii): similarly, putting $g(x) = f(e^x)$ and $g(x) = \log f(e^x)$. \square

Related:

Theorem 1.5 (Erdős?) Let $f: \mathbb{Z}^+ \rightarrow (0, \infty)$ be a function satisfying $f(mn) = f(m)f(n)$ for all $m, n \in \mathbb{Z}^+$. (There are loads of solutions: can freely choose $f(p)$ for every prime p . But ...) Suppose that either $f(1) \leq f(2) \leq \dots$ or

$$\lim_{n \rightarrow \infty} \frac{f(n+1)}{f(n)} = 1.$$

Then there exists $c \in \mathbb{R}$ such that $f(n) = n^c$ for all n .

Proof Omitted. \square

Separation of variables

When can a function of two variables be written as a product/sum of two functions of one variable? We'll do sums, but can convert to products as in Corollary 1.4.

Let X and Y be sets and

$$f: X \times Y \rightarrow \mathbb{R}$$

a function. Or can replace \mathbb{R} by any abelian group. We seek functions

$$g: X \rightarrow \mathbb{R}, \quad h: Y \rightarrow \mathbb{R}$$

such that

$$\forall x \in X, y \in Y, \quad f(x, y) = g(x) + h(y). \quad (2)$$

Basic questions:

A Are there *any* pairs of functions (g, h) satisfying (2)?

B How can we construct all such pairs?

C How many such pairs are there? Clear that if there are any, there are many, by adding/subtracting constants.

I got up to here in the first class, and was going to lecture the rest of this section in the second class, but in the end decided not to. What I actually lectured resumes at the start of Section 2. But for completeness, here's the rest of this section.

Attempt to recover g and h from f . Key insight:

$$f(x, y) - f(x_0, y) = g(x) - g(x_0)$$

($x, x_0 \in X, y \in Y$). No h s involved!

First lemma: g and h are determined by f , up to additive constant.

Lemma 1.6 Let $g: X \rightarrow \mathbb{R}$ and $h: Y \rightarrow \mathbb{R}$ be functions. Define $f: X \times Y \rightarrow \mathbb{R}$ by (2). Let $x_0 \in X$ and $y_0 \in Y$.

Then there exist $c, d \in \mathbb{R}$ such that $c + d = f(x_0, y_0)$ and

$$g(x) = f(x, y_0) - c \quad \forall x \in X, \quad (3)$$

$$h(y) = f(x_0, y) - d \quad \forall y \in Y. \quad (4)$$

Proof Put $y = y_0$ in (2): then

$$g(x) = f(x, y_0) - c \quad \forall x \in X$$

where $c = h(y_0)$. Similarly

$$h(y) = f(x_0, y) - d \quad \forall y \in Y$$

where $d = g(x_0)$. Now

$$c + d = g(x_0) + h(y_0) = f(x_0, y_0)$$

by (2). □

But given f (and x_0 and y_0), is every pair (g, h) of this form a solution of (2)? Not necessarily (but it's easy to say when)...

Lemma 1.7 *Let $f: X \times Y \rightarrow \mathbb{R}$ be a function. Let $x_0 \in X$, $y_0 \in Y$, and $c, d \in \mathbb{R}$ with $c + d = f(x_0, y_0)$. Define $g: X \rightarrow \mathbb{R}$ by (3) and $h: Y \rightarrow \mathbb{R}$ by (4). If*

$$f(x, y_0) + f(x_0, y) = f(x, y) + f(x_0, y_0) \quad \forall x \in X, y \in Y$$

then

$$f(x, y) = g(x) + h(y) \quad \forall x \in X, y \in Y.$$

Proof For all $x \in X$ and $y \in Y$,

$$g(x) + h(y) = f(x, y_0) + f(x_0, y) - c - d = f(x, y_0) + f(x_0, y) - f(x_0, y_0),$$

etc. □

Can now answer the basic questions.

Existence of decompositions (A):

Proposition 1.8 *Let $f: X \times Y \rightarrow \mathbb{R}$. TFAE:*

i. *there exist $g: X \rightarrow \mathbb{R}$ and $h: Y \rightarrow \mathbb{R}$ such that*

$$f(x, y) = g(x) + h(y) \quad \forall x \in X, y \in Y$$

ii. *$f(x, y') + f(x', y) = f(x, y) + f(x', y')$ for all x, x', y, y' .*

Proof (i) \Rightarrow (ii): trivial.

(ii) \Rightarrow (i): trivial if $X = \emptyset$ or $Y = \emptyset$. Otherwise, choose $x_0 \in X$ and $y_0 \in Y$; then use Lemma 1.7 with $c = 0$ and $d = f(x_0, y_0)$. □

Classification of decompositions (B):

Proposition 1.9 *Let $f: X \times Y \rightarrow \mathbb{R}$ be a function satisfying the equivalent conditions of Proposition 1.8, and let $x_0 \in X$ and $y_0 \in Y$. Then a pair of functions $(g: X \rightarrow \mathbb{R}, h: Y \rightarrow \mathbb{R})$ satisfies (2) if and only if there exist $c, d \in \mathbb{R}$ satisfying $c + d = f(x_0, y_0)$, (3) and (4).*

Proof Follows from Lemmas 1.6 and 1.7. □

Number of decompositions (C) (really: dim of solution-space):

Corollary 1.10 *Let $f: X \times Y \rightarrow \mathbb{R}$ with X, Y nonempty. Either there are no pairs (g, h) satisfying (2), or for any pair (g, h) satisfying (2), the set of all such pairs is the 1-dimensional space*

$$\{(g + a, h - a) : a \in \mathbb{R}\}. \quad \square$$

2 Shannon entropy

Week II (14 Feb)

Recap, including Erdős theorem. No separation of variables!

The many meanings of the word *entropy*. Ordinary entropy, relative entropy, conditional entropy, joint entropy, cross entropy; entropy on finite and infinite spaces; quantum versions; entropy in topological dynamics; ... Today we stick to the very simplest kind: Shannon entropy of a probability distribution on a finite set.

Let $\mathbf{p} = (p_1, \dots, p_n)$ be a probability distribution on $\{1, \dots, n\}$ (i.e. $p_i \geq 0$, $\sum p_i = 1$). The **(Shannon) entropy** of \mathbf{p} is

$$H(\mathbf{p}) = - \sum_{i: p_i > 0} p_i \log p_i = \sum_{i: p_i > 0} p_i \log \frac{1}{p_i}.$$

The sum is over all $i \in \{1, \dots, n\}$ such that $p_i \neq 0$; equivalently, can sum over all $i \in \{1, \dots, n\}$ but with the convention that $0 \log 0 = 0$.

Ways of thinking about entropy:

- Disorder.
- Uniformity. Will see that uniform distribution has greatest entropy among all distributions on $\{1, \dots, n\}$.
- Expected surprise. Think of $\log(1/p_i)$ as your surprise at learning that an event of probability p_i has occurred. The smaller p_i is, the more surprised you are. Then $H(\mathbf{p})$ is the expected value of the surprise: how surprised you expect to be!
- Information. Similar to expected surprise. Think of $\log(1/p_i)$ as the information that you gain by observing an event of probability p_i . The smaller p_i is, the rarer the event is, so the more remarkable it is. Then $H(\mathbf{p})$ is the average amount of information per event.
- Lack of information (!). Dual viewpoints in information theory. E.g. if \mathbf{p} represents noise, high entropy means more noise. Won't go into this.
- Genericity. In context of thermodynamics, entropy measures how generic a state a system is in. Closely related to 'lack of information'.

First properties:

- $H(\mathbf{p}) \geq 0$ for all \mathbf{p} , with equality iff $\mathbf{p} = (0, \dots, 0, 1, 0, \dots, 0)$. Least uniform distribution.
- $H(\mathbf{p}) \leq \log n$ for all \mathbf{p} , with equality iff $\mathbf{p} = (1/n, \dots, 1/n)$. Most uniform distribution. Proof that $H(\mathbf{p}) \leq \log n$ uses concavity of \log :

$$H(\mathbf{p}) = \sum_{i: p_i > 0} p_i \log\left(\frac{1}{p_i}\right) \leq \log\left(\sum_{i: p_i > 0} p_i \frac{1}{p_i}\right) \leq \log n.$$

- $H(\mathbf{p})$ is continuous in \mathbf{p} . (Uses $\lim_{x \rightarrow 0^+} x \log x = 0$.)

Remark 2.1 Base of logarithm usually taken to be e (for theory) or 2 (for examples and in information theory/digital communication). Changing base of logarithm scales H by constant factor—harmless!

Examples 2.2 Use \log_2 here.

- i. Uniform distribution on 2^k elements:

$$H\left(\frac{1}{2^k}, \dots, \frac{1}{2^k}\right) = \log_2(2^k) = k.$$

Interpretation: knowing results of k fair coin tosses gives k bits of information.

- ii.

$$\begin{aligned} H\left(\frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \frac{1}{8}\right) &= \frac{1}{2} \log_2 2 + \frac{1}{4} \log_2 4 + \frac{1}{8} \log_2 8 + \frac{1}{8} \log_2 8 \\ &= \frac{1}{2} \cdot 1 + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 + \frac{1}{8} \cdot 3 \\ &= 1\frac{3}{4}. \end{aligned}$$

Interpretation: consider a language with alphabet A, B, C, D, with frequencies $1/2, 1/4, 1/8, 1/8$. We want to send messages encoded in binary. Compare Morse code: use short code sequences for common letters. The most efficient unambiguous code encodes a letter of frequency 2^{-k} as a binary string of length k : e.g. here, could use

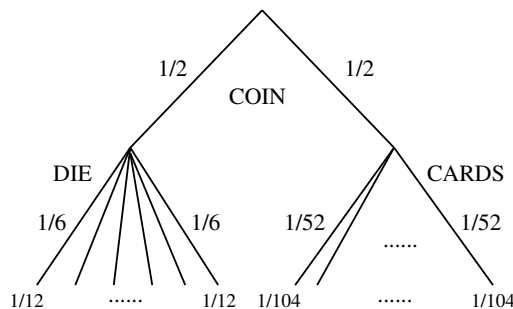
A: 0, B: 10, C: 110, D: 111.

Then code messages are unambiguous: e.g. 11010011110 can only be CBADB. Since $k = \log_2(1/p_i)$, mean number of bits per letter is then $\sum_i p_i \log(1/p_i) = H(\mathbf{p}) = 1\frac{3}{4}$.

- iii. That example was special in that all the probabilities were integer powers of 2. But... Can still make sense of this when probabilities aren't powers of 2 (Shannon's first theorem). E.g. frequency distribution $\mathbf{p} = (p_1, \dots, p_{26})$ of letters in English has $H(\mathbf{p}) \approx 4$, so can encode English in about 4 bits/letter. So, it's as if English had only 16 letters, used equally often.

Will now explain a more subtle property of entropy. Begin with example.

Example 2.3 Flip a coin. If it's heads, roll a die. If it's tails, draw from a pack of cards. So final outcome is either a number between 1 and 6 or a card. There are $6 + 52 = 58$ possible final outcomes, with probabilities as shown (assuming everything unbiased):



How much information do you expect to get from observing the outcome?

- You know result of coin flip, giving $H(1/2, 1/2) = 1$ bit of info.
- With probability $1/2$, you know result of die roll: $H(1/6, \dots, 1/6) = \log_2 6$ bits of info.
- With probability $1/2$, you know result of card draw: $H(1/52, \dots, 1/52) = \log_2 52$ bits.

In total:

$$1 + \frac{1}{2} \log_2 6 + \frac{1}{2} \log_2 52$$

bits of info. This suggests

$$H\left(\underbrace{\frac{1}{12}, \dots, \frac{1}{12}}_6, \underbrace{\frac{1}{104}, \dots, \frac{1}{104}}_{52}\right) = 1 + \frac{1}{2} \log_2 6 + \frac{1}{2} \log_2 52.$$

Can check true! Now formulate general rule.

The chain rule Write

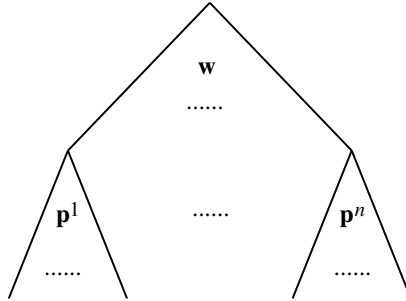
$$\Delta_n = \{\text{probability distributions on } \{1, \dots, n\}\}.$$

Geometrically, this is a simplex of dimension $n - 1$. Given

$$\mathbf{w} \in \Delta_n, \quad \mathbf{p}^1 \in \Delta_{k_1}, \dots, \mathbf{p}^n \in \Delta_{k_n},$$

get composite distribution

$$\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^n) = (w_1 p_{k_1}^1, \dots, w_1 p_{k_1}^1, \dots, w_n p_1^n, \dots, w_n p_{k_n}^n) \in \Delta_{k_1 + \dots + k_n}$$



(For cognoscenti: this defines an operad structure on the simplices.)

Easy calculation¹ proves **chain rule**:

$$H(\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^n)) = H(\mathbf{w}) + \sum_{i=1}^n w_i H(\mathbf{p}^i).$$

Special case: $\mathbf{p}^1 = \dots = \mathbf{p}^n$. For $\mathbf{w} \in \Delta_n$ and $\mathbf{p} \in \Delta_m$, write

$$\mathbf{w} \otimes \mathbf{p} = \mathbf{w} \circ \underbrace{(\mathbf{p}, \dots, \mathbf{p})}_n = (w_1 p_1, \dots, w_1 p_m, \dots, w_n p_1, \dots, w_n p_m) \in \Delta_{nm}.$$

¹This is completely straightforward, but can be made even more transparent by first observing that the function $f(x) = -x \log x$ is a ‘nonlinear derivation’, i.e. $f(xy) = xf(y) + f(x)y$. In fact, $-x \log x$ is the *only* measurable function F with this property (up to a constant factor), since if we put $g(x) = F(x)/x$ then $g(xy) = g(y) + g(x)$ and so $g(x) \propto \log x$.

This is joint probability distribution if the two things are independent. Then chain rule implies **multiplicativity**:

$$H(\mathbf{w} \otimes \mathbf{p}) = H(\mathbf{w}) + H(\mathbf{p}).$$

Interpretation: information from two independent observations is sum of information from each.

Where are the functional equations?

For each $n \geq 1$, have function $H: \Delta_n \rightarrow \mathbb{R}^+ = [0, \infty)$. Faddeev² showed:

Theorem 2.4 (Faddeev, 1956) Take functions $(I: \Delta_n \rightarrow \mathbb{R}^+)_{n \geq 1}$. TFAE:

- i. the functions I are continuous and satisfy the chain rule;
- ii. $I = cH$ for some $c \in \mathbb{R}^+$.

That is: up to a constant factor, Shannon entropy is uniquely characterized by continuity and chain rule.

Should we be disappointed to get *scalar multiples* of H , not H itself? No: recall that different scalar multiples correspond to different choices of the base for log.

Rest of this section: proof of Faddeev's theorem.

Certainly (ii) \Rightarrow (i). Now take I satisfying (i).

Write $\mathbf{u}_n = (1/n, \dots, 1/n) \in \Delta_n$. Strategy: think about the sequence $(I(\mathbf{u}_n))_{n \geq 1}$. It should be $(c \log n)_{n \geq 1}$ for some constant c .

Lemma 2.5 i. $I(\mathbf{u}_{mn}) = I(\mathbf{u}_m) + I(\mathbf{u}_n)$ for all $m, n \geq 1$.

ii. $I(\mathbf{u}_1) = 0$.

Proof For (i), $\mathbf{u}_{mn} = \mathbf{u}_m \otimes \mathbf{u}_n$, so

$$I(\mathbf{u}_{mn}) = I(\mathbf{u}_m \otimes \mathbf{u}_n) = I(\mathbf{u}_m) + I(\mathbf{u}_n)$$

(by multiplicativity). For (ii), take $m = n = 1$ in (i). □

Theorem 1.5 (Erdős) *would* now tell us that $I(\mathbf{u}_n) = c \log n$ for some constant c (putting $f(n) = \exp(I(\mathbf{u}_n))$). But to conclude that, we need one of the two alternative hypotheses of Theorem 1.5 to be satisfied. We prove the second one, on limits. This takes some effort.

Lemma 2.6 $I(1, 0) = 0$.

Proof We compute $I(1, 0, 0)$ in two ways. First,

$$I(1, 0, 0) = I((1, 0) \circ ((1, 0), \mathbf{u}_1)) = I(1, 0) + 1 \cdot I(1, 0) + 0 \cdot I(\mathbf{u}_1) = 2I(1, 0).$$

Second,

$$I(1, 0, 0) = I((1, 0) \circ (\mathbf{u}_1, \mathbf{u}_2)) = I(1, 0) + 1 \cdot I(\mathbf{u}_1) + 0 \cdot I(\mathbf{u}_2) = I(1, 0)$$

since $I(\mathbf{u}_1) = 0$. Hence $I(1, 0) = 0$. □

²Dmitry Faddeev, father of the physicist Ludvig Faddeev.

To use Erdős, need $I(\mathbf{u}_{n+1}) - I(\mathbf{u}_n) \rightarrow 0$ as $n \rightarrow \infty$. Can *nearly* prove that:

Lemma 2.7 $I(\mathbf{u}_{n+1}) - \frac{n}{n+1}I(\mathbf{u}_n) \rightarrow 0$ as $n \rightarrow \infty$.

Proof We have

$$\mathbf{u}_{n+1} = \left(\frac{n}{n+1}, \frac{1}{n}\right) \circ (\mathbf{u}_n, \mathbf{u}_1),$$

so by the chain rule and $I(\mathbf{u}_1) = 0$,

$$I(\mathbf{u}_{n+1}) = I\left(\frac{n}{n+1}, \frac{1}{n}\right) + \frac{n}{n+1}I(\mathbf{u}_n).$$

So

$$I(\mathbf{u}_{n+1}) - \frac{n}{n+1}I(\mathbf{u}_n) = I\left(\frac{n}{n+1}, \frac{1}{n+1}\right) \rightarrow I(1, 0) = 0$$

as $n \rightarrow \infty$, by continuity and Lemma 2.6. \square

To improve this to $I(\mathbf{u}_{n+1}) - I(\mathbf{u}_n) \rightarrow 0$, use a general result that has nothing to do with entropy:

Lemma 2.8 Let $(a_n)_{n \geq 1}$ be a sequence in \mathbb{R} such that $a_{n+1} - \frac{n}{n+1}a_n \rightarrow 0$ as $n \rightarrow \infty$. Then $a_{n+1} - a_n \rightarrow 0$ as $n \rightarrow \infty$.

Proof Omitted; uses Cesàro convergence.³

Although I omitted this proof in class, I'll include it here. I'll follow the argument in Feinstein, The Foundations of Information Theory, around p.7.

It is enough to prove that $a_n/(n+1) \rightarrow 0$ as $n \rightarrow \infty$. Write $b_1 = a_1$ and $b_n = a_n - \frac{n-1}{n}a_{n-1}$ for $n \geq 2$. Then $na_n = nb_n + (n-1)a_{n-1}$ for all $n \geq 2$, so

$$na_n = nb_n + (n-1)b_{n-1} + \dots + 1b_1$$

for all $n \geq 1$. Dividing through by $n(n+1)$ gives

$$\frac{a_n}{n+1} = \frac{1}{2} \cdot \text{mean}(b_1, b_2, b_2, b_3, b_3, b_3, \dots, \underbrace{b_n, \dots, b_n}_n).$$

Since $b_n \rightarrow 0$ as $n \rightarrow \infty$, the sequence

$$b_1, b_2, b_2, b_3, b_3, b_3, \dots, \underbrace{b_n, \dots, b_n}_n, \dots$$

also converges to 0. Now a general result of Cesàro states that if a sequence (x_r) converges to ℓ then the sequence (\bar{x}_r) also converges to ℓ , where $\bar{x}_r = (x_1 + \dots + x_r)/r$. Applying this to the sequence above implies that

$$\text{mean}(b_1, b_2, b_2, b_3, b_3, b_3, \dots, \underbrace{b_n, \dots, b_n}_n) \rightarrow 0 \text{ as } n \rightarrow \infty.$$

Hence $a_n/(n+1) \rightarrow 0$ as $n \rightarrow \infty$, as required. \square

We can now deduce what $I(\mathbf{u}_n)$ is:

Lemma 2.9 There exists $c \in \mathbb{R}^+$ such that $I(\mathbf{u}_n) = c \log n$ for all $n \geq 1$.

³Xǐlíng Zhāng pointed out that this is also a consequence of Stolz's lemma—or as Wikipedia calls it, the [Stolz–Cesàro theorem](#).

Proof We have $I(\mathbf{u}_{mn}) = I(\mathbf{u}_m) + I(\mathbf{u}_n)$, and by last two lemmas,

$$I(\mathbf{u}_{n+1}) - I(\mathbf{u}_n) \rightarrow 0 \text{ as } n \rightarrow \infty.$$

So can apply Erdős's theorem (1.5) with $f(n) = \exp(I(\mathbf{u}_n))$ to get $f(n) = n^c$ for some constant $c \in \mathbb{R}$. So $I(\mathbf{u}_n) = c \log n$, and $c \geq 0$ since I maps into \mathbb{R}^+ . \square

We now know that $I = cH$ on the *uniform* distributions \mathbf{u}_n . It might seem like we still have a mountain to climb to get to $I = cH$ for *all* distributions. But in fact, it's easy.

Lemma 2.10 $I(\mathbf{p}) = cH(\mathbf{p})$ whenever p_1, \dots, p_n are rational.

Proof Write

$$\mathbf{p} = \left(\frac{k_1}{k}, \dots, \frac{k_n}{k} \right)$$

where $k_1, \dots, k_n \in \mathbb{Z}$ and $k = k_1 + \dots + k_n$. Then

$$\mathbf{p} \circ (\mathbf{u}_{k_1}, \dots, \mathbf{u}_{k_n}) = \mathbf{u}_k.$$

Since I satisfies the chain rule and $I(\mathbf{u}_r) = cH(\mathbf{u}_r)$ for all r ,

$$I(\mathbf{p}) + \sum_{i=1}^n p_i \cdot cH(\mathbf{u}_{k_i}) = cH(\mathbf{u}_k).$$

But since cH also satisfies the chain rule,

$$cH(\mathbf{p}) + \sum_{i=1}^n p_i \cdot cH(\mathbf{u}_{k_i}) = cH(\mathbf{u}_k),$$

giving the result. \square

Theorem 2.4 follows by continuity.

Week III (21 Feb)

Recap of last time: Δ_n , H , chain rule.

Information is a slippery concept to reason about. One day it will seem intuitively clear that the distribution (0.5, 0.5) is 'more informative' than (0.9, 0.1), and the next day your intuition will say the opposite. So to make things concrete, it's useful to concentrate on one particular framework: coding.

Slogan:

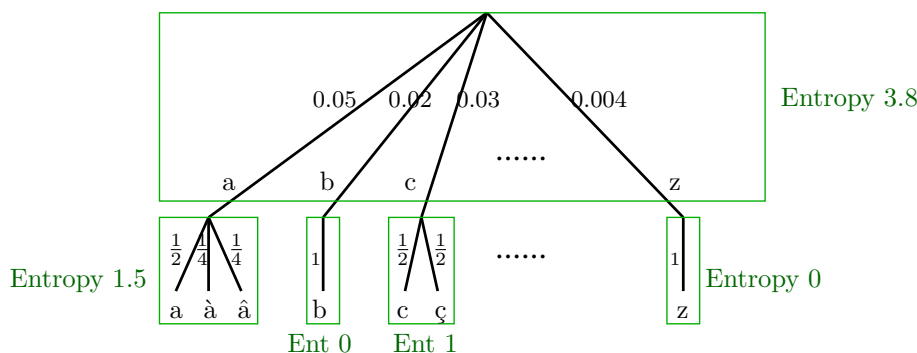
Entropy is average number of bits/symbol in an optimal encoding.

Consider language with alphabet A, B, ..., used with frequencies p_1, p_2, \dots, p_n . Want to encode efficiently in binary.

- Revisit Example 2.2(ii).
- In general: if p_1, \dots, p_n are all powers of 2, there is an unambiguous encoding where i th letter is encoded as string of length $\log_2(1/p_i)$. So mean bits/symbol = $\sum_i p_i \log_2(1/p_i) = H(\mathbf{p})$.

- Shannon's first theorem: for any $\mathbf{p} \in \Delta_n$, in any unambiguous encoding, mean bits/symbol $\geq H(\mathbf{p})$; moreover, if you're clever, can do it in $< H(\mathbf{p}) + \varepsilon$ for any $\varepsilon > 0$.
- When p_i s aren't powers of 2, do this by using *blocks* of symbols rather than individual symbols.

Chain rule: how many bits/symbol on average to encode French, including accents?



We need

$$\underbrace{3.8}_{\text{bits for actual letters}} + \underbrace{(0.05 \times 1.5 + 0.02 \times 0 + 0.03 \times 1 + \dots + 0.004 \times 0)}_{\text{bits for accents}}$$

bits/symbol. (Convention: *letters* are a, b, c, ...; *symbols* are a, à, â, b, c, ç, ...) By the chain rule, this is equal to the entropy of the composite distribution

$$(0.05 \times \frac{1}{2}, 0.05 \times \frac{1}{4}, 0.05 \times \frac{1}{4}, 0.02 \times 1, \dots, 0.004 \times 1).$$

Relative entropy

Let $\mathbf{p}, \mathbf{r} \in \Delta_n$. The **entropy of \mathbf{p} relative to \mathbf{r}** is

$$H(\mathbf{p} \parallel \mathbf{r}) = \sum_{i: p_i > 0} p_i \log\left(\frac{p_i}{r_i}\right).$$

Also called **Kullback–Leibler divergence**, **relative information**, or **information gain**.

First properties:

- $H(\mathbf{p} \parallel \mathbf{r}) \geq 0$. Not obvious, as $\log(p_i/r_i)$ is sometimes positive and sometimes negative. For since log is concave,

$$H(\mathbf{p} \parallel \mathbf{r}) = - \sum_{i: p_i > 0} p_i \log\left(\frac{r_i}{p_i}\right) \geq - \log\left(\sum_{i: p_i > 0} p_i \frac{r_i}{p_i}\right) \geq - \log 1 = 0.$$

- $H(\mathbf{p} \parallel \mathbf{r}) = 0$ if and only if $\mathbf{p} = \mathbf{r}$. Evidence so far suggests relative entropy is something like a distance. That's wrong in that it's not a metric, but it's not too terribly wrong. Will come back to this.

- $H(\mathbf{p} \parallel \mathbf{r})$ can be arbitrarily large (even for fixed n). E.g.

$$H((1/2, 1/2) \parallel (t, 1-t)) \rightarrow \infty \text{ as } t \rightarrow \infty,$$

and in fact $H(\mathbf{p} \parallel \mathbf{r}) = \infty$ if $p_i > 0 = r_i$ for some i .

- Write

$$\mathbf{u}_n = (1/n, \dots, 1/n) \in \Delta_n.$$

Then

$$H(\mathbf{p} \parallel \mathbf{u}_n) = \log n - H(\mathbf{p}).$$

So entropy is pretty much a special case of relative entropy.

- $H(\mathbf{p} \parallel \mathbf{r}) \neq H(\mathbf{r} \parallel \mathbf{p})$. E.g.

$$H(\mathbf{u}_2 \parallel (0, 1)) = \infty,$$

$$H((0, 1) \parallel \mathbf{u}_2) = \log 2 - H((0, 1)) = \log 2.$$

Will come back to this too.

Coding interpretation Convenient fiction: for each ‘language’ \mathbf{p} , there is an encoding for \mathbf{p} using $\log(1/p_i)$ bits for the i th symbol, hence with mean bits/symbol = $H(\mathbf{p})$ exactly. Call this ‘machine \mathbf{p} ’.

We have

$$\begin{aligned} H(\mathbf{p} \parallel \mathbf{r}) &= \sum p_i \log\left(\frac{1}{r_i}\right) - \sum p_i \log\left(\frac{1}{p_i}\right) \\ &= (\text{bits/symbol to encode language } \mathbf{p} \text{ using machine } \mathbf{r}) \\ &\quad - (\text{bits/symbol to encode language } \mathbf{p} \text{ using machine } \mathbf{p}) \end{aligned}$$

So relative entropy is the number of extra bits needed if you use the wrong machine. Or: penalty you pay for using the wrong machine. Explains why $H(\mathbf{p} \parallel \mathbf{r}) \geq 0$ with equality if $\mathbf{p} = \mathbf{r}$.

If $r_i = 0$ then in machine \mathbf{r} , the i th symbol has an infinitely long code word. Or if you like: if $r_i = 2^{-1000}$ then its code word has length 1000. So if also $p_i > 0$ then for language \mathbf{p} encoded using machine \mathbf{r} , average bits/symbol = ∞ . This explains why $H(\mathbf{p} \parallel \mathbf{r}) = \infty$.

Taking $\mathbf{r} = \mathbf{u}_n$,

$$H(\mathbf{p} \parallel \mathbf{u}_n) = \log n - H(\mathbf{p}) \leq \log n.$$

Explanation: in machine \mathbf{u}_n , every symbol is encoded with $\log n$ bits, so the average extra bits/symbol caused by using machine \mathbf{u}_n instead of machine \mathbf{p} is $\leq \log n$.

Now a couple of slightly more esoteric comments, pointing in different mathematical directions (both away from functional equations). Tune out if you want...

Measure-theoretic perspective Slogan:

All entropy is relative.

Attempt to generalize definition of entropy from probability measures on finite sets to arbitrary probability measures μ : want to say $H(\mu) = -\int \log(\mu) d\mu$, but this makes no sense!

Note that the finite definition $H(p) = -\sum p_i \log p_i$ implicitly refers to counting measure...

However, can generalize *relative* entropy. Given measures μ and ν on measurable space X , define

$$H(\mu \parallel \nu) = \int_X \log\left(\frac{d\mu}{d\nu}\right) d\mu$$

where $\frac{d\mu}{d\nu}$ is Radon–Nikodym derivative. *This makes sense and is the right definition.*

People do talk about the entropy of probability distributions on \mathbb{R}^n . For instance, the entropy of a probability density function f on \mathbb{R} is usually defined as $H(f) = -\int_{\mathbb{R}} f(x) \log f(x) dx$, and it's an important result that among all density functions on \mathbb{R} with a given mean and variance, the one with the maximal entropy is the normal distribution. (This is related to the central limit theorem.) But here we're implicitly using Lebesgue measure λ on \mathbb{R} ; so there are two measures in play, λ and $f\lambda$, and $H(f) = H(f\lambda \parallel \lambda)$.

Local behaviour of relative entropy Take two close-together distributions $\mathbf{p}, \mathbf{p} + \boldsymbol{\delta} \in \Delta_n$. (So $\boldsymbol{\delta} = (\delta_1, \dots, \delta_n)$ with $\sum \delta_i = 0$.) Taylor expansion gives

$$H(\mathbf{p} + \boldsymbol{\delta} \parallel \mathbf{p}) \approx \frac{1}{2} \sum \frac{1}{p_i} \delta_i^2$$

for $\boldsymbol{\delta}$ small. Precisely:

$$H(\mathbf{p} + \boldsymbol{\delta} \parallel \mathbf{p}) = \frac{1}{2} \sum \frac{1}{p_i} \delta_i^2 + O(\|\boldsymbol{\delta}\|^3) \text{ as } \boldsymbol{\delta} \rightarrow \mathbf{0}.$$

(Here $\|\cdot\|$ is any norm on \mathbb{R}^n . It doesn't matter which, as they're all equivalent.) So:

Locally, $H(- \parallel -)$ is like a squared distance.

In particular, locally (to second order) it's symmetric.

The square root of relative entropy is *not* a metric on Δ_n : not symmetric and fails triangle inequality. (E.g. put $\mathbf{p} = (0.9, 0.1)$, $\mathbf{q} = (0.2, 0.8)$, $\mathbf{r} = (0.1, 0.9)$. Then $\sqrt{H(\mathbf{p} \parallel \mathbf{q})} + \sqrt{H(\mathbf{q} \parallel \mathbf{r})} < \sqrt{H(\mathbf{p} \parallel \mathbf{r})}$.) But using it as a 'local distance' leads to important things, e.g. Fisher information (statistics), the Jeffreys prior (Bayesian statistics), and the whole subject of information geometry.

Next week: a unique characterization of relative entropy.

Recap: entropy as mean bits/symbol in optimal encoding; definition of relative entropy; relative entropy as extra cost of using wrong machine.

Write

$$A_n = \{(\mathbf{p}, \mathbf{r}) \in \Delta_n \times \Delta_n : r_i = 0 \implies p_i = 0\}.$$

So $H(\mathbf{p} \parallel \mathbf{r}) < \infty \iff (\mathbf{p}, \mathbf{r}) \in A_n$. (Secretly, A stands for ‘absolutely continuous’.) Then for each $n \geq 1$, we have the function

$$H(- \parallel -): A_n \rightarrow \mathbb{R}^+.$$

Properties:

Measurability $H(- \parallel -)$ is measurable. If you don’t know what that means, ignore: *very* mild condition. Every function that anyone has ever written down a formula for, or ever will, is measurable.

Permutation-invariance E.g. $H((p_1, p_2, p_3) \parallel (r_1, r_2, r_3)) = H((p_2, p_3, p_1) \parallel (r_2, r_3, r_1))$. It’s *that* kind of symmetry, not the kind where you swap \mathbf{p} and \mathbf{r} .

Vanishing $H(\mathbf{p} \parallel \mathbf{p}) = 0$ for all \mathbf{p} . Remember: $H(\mathbf{p} \parallel \mathbf{r})$ is *extra* cost of using machine \mathbf{r} (instead of machine \mathbf{p}) to encode language \mathbf{p} .

Chain rule For all

$$(\mathbf{w}, \tilde{\mathbf{w}}) \in A_n, (\mathbf{p}^1, \tilde{\mathbf{p}}^1) \in A_{k_1}, \dots, (\mathbf{p}^n, \tilde{\mathbf{p}}^n) \in A_{k_n},$$

we have

$$H(\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^n) \parallel \tilde{\mathbf{w}} \circ (\tilde{\mathbf{p}}^1, \dots, \tilde{\mathbf{p}}^n)) = H(\mathbf{w} \parallel \tilde{\mathbf{w}}) + \sum_{i=1}^n w_i H(\mathbf{p}^i \parallel \tilde{\mathbf{p}}^i).$$

To understand chain rule for relative entropy:

Example 2.11 Recall the letter/accent tree from last time.

Swiss French and Canadian French are written using the same alphabet and accents, but with slightly different words, hence different frequencies of letters and accents.

Consider Swiss French and Canadian French:

$$\begin{aligned} \mathbf{w} &\in \Delta_{26} : \text{frequencies of letters in Swiss} \\ \tilde{\mathbf{w}} &\in \Delta_{26} : \text{frequencies of letters in Canadian} \end{aligned}$$

and then

$$\begin{aligned} \mathbf{p}^1 &\in \Delta_3 : \text{frequencies of accents on ‘a’ in Swiss} \\ \tilde{\mathbf{p}}^1 &\in \Delta_3 : \text{frequencies of accents on ‘a’ in Canadian} \\ &\vdots \\ \mathbf{p}^{26} &\in \Delta_1 : \text{frequencies of accents on ‘z’ in Swiss} \\ \tilde{\mathbf{p}}^{26} &\in \Delta_1 : \text{frequencies of accents on ‘z’ in Canadian.} \end{aligned}$$

So

$\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^{26}) = \text{frequency distribution of all symbols in Swiss}$

$\tilde{\mathbf{w}} \circ (\tilde{\mathbf{p}}^1, \dots, \tilde{\mathbf{p}}^{26}) = \text{frequency distribution of all symbols in Canadian}$

where ‘symbol’ means a letter with (or without) an accent.

Now encode Swiss using Canadian machine. How much extra does it cost (in mean bits/symbol) compared to encoding Swiss using Swiss machine?

mean extra cost per symbol =

mean extra cost per letter + mean extra cost per accent.

And that’s the chain rule. Coefficient in sum is w_i , not \tilde{w}_i , because it’s Swiss we’re encoding.

Clearly any scalar multiple of relative entropy also has these four properties (measurability, permutation-invariance, vanishing, chain rule).

Theorem 2.12 Take functions $(I(- \| -): A_n \rightarrow \mathbb{R}^+)_{n \geq 1}$. TFAE:

i. the functions I satisfy measurability, permutation-invariance, vanishing and the chain rule;

ii. $I(- \| -) = cH(- \| -)$ for some $c \in \mathbb{R}^+$.

It’s hard to believe that Theorem 2.12 is new; it could have been proved in the 1950s. However, I haven’t found it in the literature. The proof that follows is inspired by work of Baez and Fritz, who in turn built on and corrected work of Petz. Cf. Baez and Fritz ([arXiv:1402.3067](https://arxiv.org/abs/1402.3067)) and Petz (cited by B&F).

Take $I(- \| -)$ satisfying (i). Define $L: (0, 1] \rightarrow \mathbb{R}^+$ by

$$L(\alpha) = I((1, 0) \| (\alpha, 1 - \alpha)).$$

Note $((1, 0), (\alpha, 1 - \alpha)) \in A_2$, so RHS is defined. If $I = H$ then $L(\alpha) = -\log \alpha$.

Lemma 2.13 Let $(\mathbf{p}, \mathbf{r}) \in A_n$ with $p_{k+1} = \dots = p_n = 0$, where $1 \leq k \leq n$. Then $r_1 + \dots + r_k > 0$, and

$$I(\mathbf{p} \| \mathbf{r}) = L(r_1 + \dots + r_k) + I(\mathbf{p}' \| \mathbf{r}')$$

where

$$\mathbf{p}' = (p_1, \dots, p_k), \quad \mathbf{r}' = \frac{(r_1, \dots, r_k)}{r_1 + \dots + r_k}.$$

Proof Case $k = n$ trivial; suppose $k < n$.

Since \mathbf{p} is a probability distribution with $p_i = 0$ for all $i > k$, there is some $i \leq k$ such that $p_i > 0$, and then $r_i > 0$ since $(\mathbf{p}, \mathbf{r}) \in A_n$. Hence $r_1 + \dots + r_k > 0$. By definition of operadic composition,

$$I(\mathbf{p} \| \mathbf{r}) = I\left((1, 0) \circ (\mathbf{p}', \mathbf{r}'') \left\| (r_1 + \dots + r_k, r_{k+1} + \dots + r_n) \circ (\mathbf{r}', \mathbf{r}'')\right.\right)$$

where \mathbf{r}'' is the normalization of (r_{k+1}, \dots, r_n) if $r_{k+1} + \dots + r_n > 0$, or is chosen arbitrarily in Δ_{n-k} otherwise. (The set Δ_{n-k} is nonempty since $k < n$.) By chain rule, this is equal to

$$L(r_1 + \dots + r_k) + 1 \cdot I(\mathbf{p}' \| \mathbf{r}') + 0 \cdot I(\mathbf{r}'' \| \mathbf{r}''),$$

and result follows. In order to use the chain rule, we needed to know that various pairs were in A_2 , A_k , etc.; that’s easily checked. \square

Lemma 2.14 $L(\alpha\beta) = L(\alpha) + L(\beta)$ for all $\alpha, \beta \in (0, 1]$.

Proof Consider

$$x := I((1, 0, 0) \parallel (\alpha\beta, \alpha(1 - \beta), 1 - \alpha)).$$

On one hand, Lemma 2.13 with $k = 1$ gives

$$x = L(\alpha\beta) + I((1) \parallel (1)) = L(\alpha\beta)$$

On other, Lemma 2.13 with $k = 2$ gives

$$x = L(\alpha) + I((1, 0) \parallel (\beta, 1 - \beta)) = L(\alpha) + L(\beta).$$

Result follows. \square

Lemma 2.15 *There is a unique constant $c \in \mathbb{R}^+$ such that $L(\alpha) = -c \log \alpha$ for all $\alpha \in (0, 1]$.*

Proof Follows from Lemma 2.14 and measurability, as in Cor 1.4. That corollary was stated under the hypothesis of continuity, but measurability would have been enough. \square

Now we come to a clever part of Baez and Fritz's argument.

Lemma 2.16 *Let $(\mathbf{p}, \mathbf{r}) \in A_n$ and suppose that $p_i > 0$ for all i . Then $I(\mathbf{p} \parallel \mathbf{r}) = cH(\mathbf{p} \parallel \mathbf{r})$.*

Proof We have $(\mathbf{p}, \mathbf{r}) \in A_n$, so $r_i > 0$ for all i . So can choose $\alpha \in (0, 1]$ such that $r_i - \alpha p_i \geq 0$ for all i .

We will compute the (well-defined) number

$$x := I\left((p_1, \dots, p_n, \underbrace{0, \dots, 0}_n) \parallel (\alpha p_1, \dots, \alpha p_n, r_1 - \alpha p_1, \dots, r_n - \alpha p_n)\right)$$

in two ways. First, by Lemma 2.13 and the vanishing property,

$$x = L(\alpha) + I(\mathbf{p} \parallel \mathbf{p}) = -c \log \alpha.$$

Second, by symmetry and then the chain rule,

$$\begin{aligned} x &= I((p_1, 0, \dots, p_n, 0) \parallel (\alpha p_1, r_1 - \alpha p_1, \dots, p_n, r_n - \alpha p_n)) \\ &= I\left(\mathbf{p} \circ ((1, 0), \dots, (1, 0)) \parallel \mathbf{r} \circ \left(\left(\alpha \frac{p_1}{r_1}, 1 - \alpha \frac{p_1}{r_1}\right), \dots, \left(\alpha \frac{p_n}{r_n}, 1 - \alpha \frac{p_n}{r_n}\right)\right)\right) \\ &= I(\mathbf{p} \parallel \mathbf{r}) + \sum_{i=1}^n p_i L\left(\alpha \frac{p_i}{r_i}\right) \\ &= I(\mathbf{p} \parallel \mathbf{r}) - c \log \alpha - cH(\mathbf{p} \parallel \mathbf{r}). \end{aligned}$$

Comparing the two expressions for x gives the result. \square

Proof of Theorem 2.12 Let $(\mathbf{p}, \mathbf{r}) \in A_n$. By symmetry, can assume

$$p_1 > 0, \dots, p_k > 0, p_{k+1} = 0, \dots, p_n = 0$$

where $1 \leq k \leq n$. Writing $R = r_1 + \dots + r_k$,

$$\begin{aligned} I(\mathbf{p} \parallel \mathbf{r}) &= L(R) + I((p_1, \dots, p_k) \parallel \frac{1}{R}(r_1, \dots, r_k)) && \text{by Lemma 2.13} \\ &= -c \log R + cH((p_1, \dots, p_k) \parallel \frac{1}{R}(r_1, \dots, r_k)) && \text{by Lemmas 2.15 and 2.16.} \end{aligned}$$

This holds for all $I(- \parallel -)$ satisfying the four conditions; in particular, it holds for $cH(- \parallel -)$. Hence $I(\mathbf{p} \parallel \mathbf{r}) = cH(\mathbf{p} \parallel \mathbf{r})$. \square

Remark 2.17 Assuming permutation-invariance, the chain rule is equivalent to a very special case:

$$\begin{aligned} H((tw_1, (1-t)w_1, w_2, \dots, w_n) \parallel (\tilde{t}\tilde{w}_1, (1-\tilde{t})\tilde{w}_1, \tilde{w}_2, \dots, \tilde{w}_n)) \\ = H(\mathbf{w} \parallel \tilde{\mathbf{w}}) + w_1 H((t, 1-t) \parallel (\tilde{t}, 1-\tilde{t})) \end{aligned}$$

for all $(\mathbf{w}, \tilde{\mathbf{w}}) \in A_n$ and $((t, 1-t), (\tilde{t}, 1-\tilde{t})) \in A_2$. (Proof: induction.)

Next week: I'll introduce a family of 'deformations' or 'quantum versions' of entropy and relative entropy. This turns out to be important if we want a balanced perspective on what biodiversity is.

Week V (7 Mar)

Recap: formulas for entropy and cross entropy (in both $\log(1/\text{something})$ and $-\log(\text{something})$ forms); chain rules for both.

Remark 2.18 I need to make explicit something we've been using implicitly for a while.

Let $\mathbf{p} \in \Delta_n$ and $\mathbf{r} \in \Delta_m$. Then we get a new distribution

$$\begin{aligned} \mathbf{p} \otimes \mathbf{r} &= (p_1 r_1, \dots, p_1 r_m, \\ &\quad \vdots \\ &\quad p_n r_1, \dots, p_n r_m) \\ &= \mathbf{p} \circ \underbrace{(\mathbf{r}, \dots, \mathbf{r})}_n \\ &\in \Delta_{nm}. \end{aligned}$$

Probabilistically, this is the joint distribution of independent random variables distributed according to \mathbf{p} and \mathbf{r} .

A special case of the chain rule for entropy:

$$H(\mathbf{p} \otimes \mathbf{r}) = H(\mathbf{p}) + H(\mathbf{r})$$

and similarly for relative entropy:

$$H(\mathbf{p} \otimes \mathbf{r} \parallel \tilde{\mathbf{p}} \otimes \tilde{\mathbf{r}}) = H(\mathbf{p} \parallel \tilde{\mathbf{p}}) + H(\mathbf{r} \parallel \tilde{\mathbf{r}}).$$

(So H and $H(- \parallel -)$ are log-like; like a higher version of Cauchy's functional equation.)

3 Deformed entropies

Shannon entropy is just a single member of a one-parameter family of entropies. In fact, there are *two* different one-parameter families of entropies, both containing Shannon entropy as a member. In some sense, these two families of entropies are equivalent, but they have different flavours.

I'll talk about both families: surprise entropies in this section, then Rényi entropies when we come to measures of diversity later on.

Definition 3.1 Let $q \in [0, \infty)$. The q -**logarithm** $\ln_q: (0, \infty) \rightarrow \mathbb{R}$ is defined by

$$\ln_q(x) = \begin{cases} \frac{x^{1-q} - 1}{1-q} & \text{if } q \neq 1, \\ \ln(x) & \text{if } q = 1. \end{cases}$$

Notation: log vs. ln.

Then $\ln_q(x)$ is continuous in q (proof: l'Hôpital's rule). Warning:

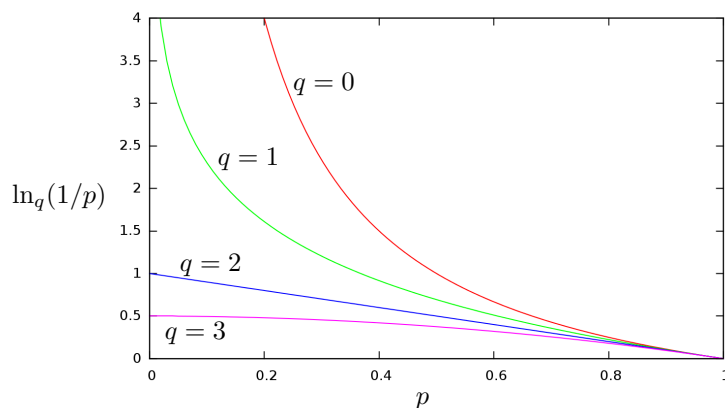
$$\ln_q(xy) \neq \ln_q(x) + \ln_q(y), \quad \ln_q(1/x) \neq -\ln_q(x).$$

First one inevitable: we already showed that scalar multiples of *actual* logarithm are the only continuous functions that convert multiplication into addition. In fact, there's quite a neat formula for $\ln_q(xy)$ in terms of $\ln_q(x)$, $\ln_q(y)$ and q : exercise!

For a probability $p \in [0, 1]$, can view

$$\ln_q(1/p)$$

as one's 'surprise' at witnessing an event with probability p . (Decreasing in p ; takes value 0 at $p = 1$.)



Definition 3.2 Let $q \in [0, \infty)$. The **surprise entropy of order q** of $\mathbf{p} \in \Delta_n$ is

$$S_q(\mathbf{p}) = \sum_{i: p_i > 0} p_i \ln_q(1/p_i) = \begin{cases} \frac{1}{1-q} (\sum p_i^q - 1) & \text{if } q \neq 1, \\ \sum p_i \ln(1/p_i) & \text{if } q = 1. \end{cases}$$

Interpretation: $S_q(\mathbf{p})$ is the *expected surprise* of an event drawn from \mathbf{p} .

Usually called 'Tsallis entropies', because Tsallis discovered them in physics in 1988, after Havrda and Charvat discovered and developed them in information

theory in 1967, and after Vajda (1968), Daróczy (1970) and Sharma and Mittal (1975) developed them further, and after Patil and Taillie used them as measures of biodiversity in 1982.

Tsallis says the parameter q can be interpreted as the ‘degree of non-extensivity’. When we come to diversity measures, I’ll give another interpretation.

Special cases:

- $S_0(\mathbf{p}) = |\{i : p_i > 0\}| - 1$
- $S_2(\mathbf{p}) = 1 - \sum_i p_i^2$ = probability that two random elements of $\{1, \dots, n\}$ (chosen according to \mathbf{p}) are different.

Properties of S_q (for fixed q):

- Permutation-invariance, e.g. $S_q(p_2, p_3, p_1) = S_q(p_1, p_2, p_3)$.
- q -chain rule:

$$S_q(\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^n)) = S_q(\mathbf{w}) + \sum_{i: w_i > 0} w_i^q S_q(\mathbf{p}^i).$$

- q -multiplicativity: as special case of chain rule,

$$S_q(\mathbf{w} \otimes \mathbf{p}) = S_q(\mathbf{w}) + \left(\sum_i w_i^q \right) S_q(\mathbf{p}).$$

Theorem 3.3 Let $q \in [0, \infty) \setminus \{1\}$. Let $(I: \Delta_n \rightarrow \mathbb{R}^+)_{n \geq 1}$ be functions. TFAE:

- I is permutation-invariant and q -multiplicative in sense above;
- $I = cS_q$ for some $c \in \mathbb{R}^+$.

No regularity condition needed! And don’t need full chain rule—just a special case, multiplicativity.

Following proof is extracted from Aczél and Daróczy’s book (Theorem 6.3.9), but they make it look way more complicated. The key point is that the multiplicativity property is not symmetric.

Proof (ii) \Rightarrow (i) easy. Assume (i). By permutation-invariance,

$$I(\mathbf{p} \otimes \mathbf{r}) = I(\mathbf{r} \otimes \mathbf{p})$$

for all $\mathbf{p} \in \Delta_n$ and $\mathbf{r} \in \Delta_m$. So by q -multiplicativity,

$$I(\mathbf{p}) + \left(\sum_i p_i^q \right) I(\mathbf{r}) = I(\mathbf{r}) + \left(\sum_i r_i^q \right) I(\mathbf{p}),$$

hence

$$\left(1 - \sum_i r_i^q \right) I(\mathbf{p}) = \left(1 - \sum_i p_i^q \right) I(\mathbf{r}).$$

Now want to get the \mathbf{p} s on one side and the \mathbf{r} s on the other, to deduce that $I(\mathbf{p})$ is proportional to $1 - \sum_i p_i^q$. But need to be careful about division by zero. Take $\mathbf{r} = \mathbf{u}_2 = (1/2, 1/2)$: then

$$\left(1 - 2^{1-q} \right) I(\mathbf{p}) = \left(1 - \sum_i p_i^q \right) I(\mathbf{u}_2)$$

for all \mathbf{p} . But $q \neq 1$, so $1 - 2^{1-q} \neq 0$. So $I = cS_q$ where $c = I(\mathbf{u}_2) \frac{1-q}{2^{1-q}-1}$. \square

Relative entropy generalizes easily too, i.e. extends all along the family from the point $q = 1$. Only ticklish point is $\ln_q(1/x)$ vs. $-\ln_q(x)$.

Definition 3.4 Let $q \in [0, \infty)$. For $\mathbf{p}, \mathbf{r} \in \Delta_n$, the **relative surprise entropy of order q** is

$$S_q(\mathbf{p} \parallel \mathbf{r}) = - \sum_{i: p_i > 0} p_i \ln_q \left(\frac{r_i}{p_i} \right) = \begin{cases} \frac{1}{1-q} \left(1 - \sum p_i^q r_i^{1-q} \right) & \text{if } q \neq 1, \\ \sum p_i \log \left(\frac{p_i}{r_i} \right) & \text{if } q = 1. \end{cases}$$

Properties:

- Permutation-invariance (as in case $q = 1$).
- q -chain rule:

$$S_q(\mathbf{w} \circ (\mathbf{p}^1, \dots, \mathbf{p}^n) \parallel \tilde{\mathbf{w}} \circ (\tilde{\mathbf{p}}^1, \dots, \tilde{\mathbf{p}}^n)) = S_q(\mathbf{w} \parallel \tilde{\mathbf{w}}) + \sum_{i: w_i > 0} w_i^q \tilde{w}_i^{1-q} S_q(\mathbf{p}^i \parallel \tilde{\mathbf{p}}^i).$$

- q -multiplicativity: as special case of chain rule,

$$S_q(\mathbf{w} \otimes \mathbf{p} \parallel \tilde{\mathbf{w}} \otimes \tilde{\mathbf{p}}) = S_q(\mathbf{w} \parallel \tilde{\mathbf{w}}) + \left(\sum_{i: w_i > 0} w_i^q \tilde{w}_i^{1-q} \right) S_q(\mathbf{p} \parallel \tilde{\mathbf{p}}).$$

Again, there's a ludicrously simple characterization theorem that needs no regularity condition. Nor does it need the vanishing condition of Theorem 2.12.

Recall notation:

$$A_n = \{(\mathbf{p}, \mathbf{r}) \in \Delta_n \times \Delta_n : r_i = 0 \implies p_i = 0\}.$$

Then $S_q(\mathbf{p} \parallel \mathbf{r}) < \infty \iff (\mathbf{p}, \mathbf{r}) \in A_n$.

Theorem 3.5 Let $q \in [0, \infty) \setminus \{1\}$. Let $(I(- \parallel -): A_n \rightarrow \mathbb{R}^+)_{n \geq 1}$ be functions. TFAE:

- $I(- \parallel -)$ is permutation-invariant and q -multiplicative in sense above;
- $I(- \parallel -) = cS_q(- \parallel -)$ for some $c \in \mathbb{R}^+$.

Proof (ii) \implies (i) easy. Assume (i). By permutation-invariance,

$$I(\mathbf{p} \otimes \mathbf{r} \parallel \tilde{\mathbf{p}} \otimes \tilde{\mathbf{r}}) = I(\mathbf{r} \otimes \mathbf{p} \parallel \tilde{\mathbf{r}} \otimes \tilde{\mathbf{p}})$$

for all $(\mathbf{p}, \tilde{\mathbf{p}}) \in A_n$ and $(\mathbf{r}, \tilde{\mathbf{r}}) \in A_m$. So by q -multiplicativity,

$$I(\mathbf{p} \parallel \tilde{\mathbf{p}}) + \left(\sum p_i^q \tilde{p}_i^{1-q} \right) I(\mathbf{r} \parallel \tilde{\mathbf{r}}) = I(\mathbf{r} \parallel \tilde{\mathbf{r}}) + \left(\sum r_i^q \tilde{r}_i^{1-q} \right) I(\mathbf{p} \parallel \tilde{\mathbf{p}}),$$

hence

$$\left(1 - \sum r_i^q \tilde{r}_i^{1-q} \right) I(\mathbf{p} \parallel \tilde{\mathbf{p}}) = \left(1 - \sum p_i^q \tilde{p}_i^{1-q} \right) I(\mathbf{r} \parallel \tilde{\mathbf{r}}).$$

Take $\mathbf{r} = (1, 0)$ and $\tilde{\mathbf{r}} = \mathbf{u}_2$: then

$$(1 - 2^{q-1}) I(\mathbf{p} \parallel \tilde{\mathbf{p}}) = I((1, 0) \parallel \mathbf{u}_2) (1 - \sum p_i^q \tilde{p}_i^{1-q})$$

for all $(\mathbf{p}, \tilde{\mathbf{p}}) \in A_n$. But $q \neq 1$, so $1 - 2^{q-1} \neq 0$. So $I(- \parallel -) = cS_q(- \parallel -)$ where $c = I((1, 0) \parallel \mathbf{u}_2) \frac{1-q}{2^{1-q}-1}$. \square

Next week: how to use probability theory to solve functional equations.