

# Solution Techniques for Large-Scale Financial Planning Problems

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# Scope of this talk

Multi-period financial planning problem

Importance of problem formulation

Exploiting structure and parallelism

Warm-start for stochastic programming problems

# Financial planning problems

Why:

- ▶ Well studied area
- ▶ Useful application
- ▶ Possible to generate large-scale problems

Stochastic programming framework:

- ▶ Multi-period structure
- ▶ Uncertain returns

# Multi-period financial planning problem

- ▶ A set of assets  $\mathcal{J} = \{1, \dots, J\}$  is given.
- ▶ Initial investment  $b$ .
- ▶ At every stage  $t = 0, \dots, T-1$  we can buy or sell any assets.
- ▶ The return of asset  $j$  at stage  $t$  is **uncertain**.

Competing objectives:

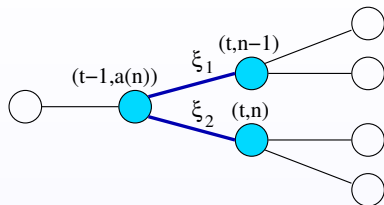
- ▶ maximize the final wealth
- ▶ minimize the associated risk

Mean-Variance formulation (Markowitz):  $\max \mathbf{E}(X) - \rho \text{Var}(X)$ .

$X$  value of the final portfolio

$\rho$  investor's attitude to risk

## Modelling with event tree



With asset  $j \in \mathcal{J}$  at node  $i = (t, n)$  we associate:

$x_{i,j}^h$  position in asset  $j$  at node  $i$

$x_{i,j}^b, x_{i,j}^s$  amount of asset  $j$  bought/sold at node  $i$

$v_j$  value of asset  $j$

$r_{j,t}$  return of asset  $j$  when held at time  $t$

$L_i, C_i$  liabilities/cash contributions at node  $i$

# Asset and Liability Management Problem I

Objective:

$$E(X) = (1 - c_t) \sum_{i \in L_T} p_i \sum_j v_j x_{i,j}^h = y$$

$$\text{Var}(X) = \sum_{i \in L_T} p_i (1 - c_t)^2 \left[ \sum_j v_j x_{i,j}^h \right]^2 - y^2$$

Constraints at each node  $i$ :

$$x_{i,j}^h = (1 + r_{i,j}) x_{a(i),j}^h + x_{i,j}^b - x_{i,j}^s \quad (\text{inventory})$$

$$\sum_j (1 + c_t) v_j x_{i,j}^b + L_i = \sum_j (1 - c_t) v_j x_{i,j}^s + C_i \quad (\text{cash balance})$$

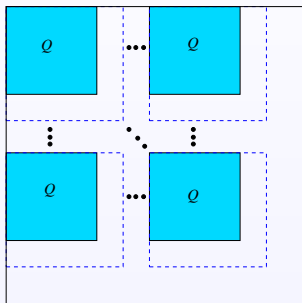
## Asset and Liability Management Problem II

$$\begin{aligned} \max_{x,y \geq 0} \quad & y - \rho \left[ \sum_{i \in L_T} p_i [(1 - c_t) \sum_j v_j x_{i,j}^h]^2 - y^2 \right] \\ \text{s.t.} \quad & (1 - c_t) \sum_{i \in L_T} p_i \sum_j v_j x_{i,j}^h = y \\ & (1 + r_{i,j}) x_{a(i),j}^h = x_{i,j}^h - x_{i,j}^b + x_{i,j}^s \quad \forall i, \forall j \\ & \sum_j (1 + c_t) v_j x_{i,j}^b + L_i = \sum_j (1 - c_t) v_j x_{i,j}^s + C_i \quad \forall i \\ & \sum_j (1 + c_t) v_j x_{0,j}^b = b \end{aligned}$$

## Structure of the objective I

Straightforward representation:

$$\begin{aligned} E(X) - \rho \text{Var}(X) &= E(X) - \rho [E(X^2) - E(X)^2] \\ &= \sum_{i \in L_T} p_i \sum_j v_j x_{ij}^h - \rho \left[ \sum_{i \in L_T} p_i \sum_j (v_j x_{ij}^h)^2 - \left[ \sum_{i \in L_T} p_i \sum_j v_j x_{ij}^h \right]^2 \right] \end{aligned}$$



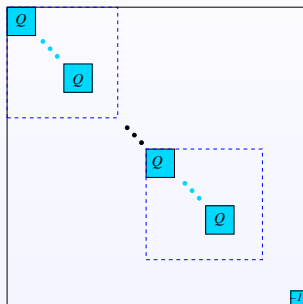
Dense, positive semidefinite Hessian

## Structure of the objective II

Alternative representation:

$$E(X) - \rho \text{Var}(X) = y - \rho \left[ \sum_{i \in L_T} p_i \sum_j (v_j x_{ij}^h)^2 - y^2 \right]$$

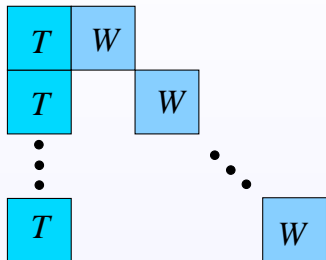
$$\text{where: } y = \sum_{i \in L_T} p_i \sum_j v_j x_{ij}^h$$



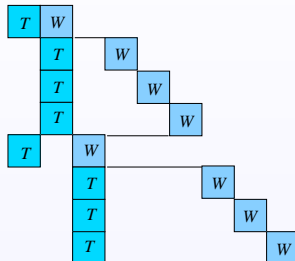
Sparse, indefinite Hessian

# Structure in the constraint matrix

Stochastic programming problems give rise to matrices with **block-angular structure**:



$$T_i x^1 + W_i y_i = b_i$$



$$T_{l_t} x_{a(l_t)} + W_{l_t} x_{l_t} = b_{l_t}$$

# The curse of dimensions

The deterministic equivalent formulation produces problems of extremely large size, even when starting from a small core.

Example: <code>fxm</code>	rows	cols	nonzeros
Core matrix	330	457	2,566
3 stages, 6 nodes:	6,200	9,492	54,589
4 stages, 16 nodes:	386,940	517,282	4,518,039

- ▶ A detailed description produces robust decisions
- ▶ Detailed event trees can be very large
- ▶ The dimensions involved explode

However, remember the presence of [structure](#)!

# The way forward

Enter **interior point methods**:

- ▶ IPM solvers are available in the community (CPLEX Barrier, PCx, HOPDM, etc.)
- ▶ Competitiveness of IPMs grows with the problem size
- ▶ Parallel implementations are possible

And we can exploit the structure:

- ▶ Linear algebra: structure-exploiting parallel software **OOPS**
- ▶ Algorithmically: **warm-start** for stochastic problems in IPMs

# OOPS - Object Oriented Parallel (Interior Point) Solver

Key advantages of exploiting the structure in the problem:

- ▶ Faster linear algebra
- ▶ Reduced memory use (by use of implicit factorization)
- ▶ Possibility to exploit (massive) parallelism
- ▶ We assume that structure is known!

OOPS is a general purpose (parallel) Interior Point solver

- ▶ Not tuned to any particular hardware or problem
- ▶ OOPS currently solves LP/QP problems
- ▶ NLP extension solves nonlinear financial planning problems

# Performance of OOPS

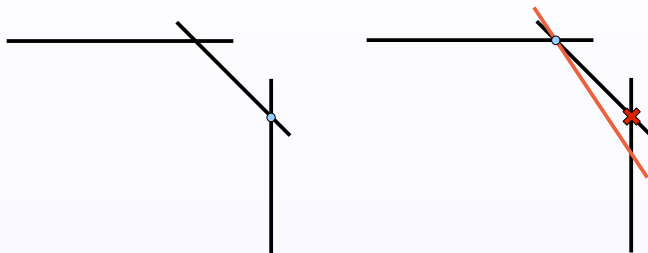
Problem	Stgs	Blks	Assets	Scens	Cons	Vars	iter	time	procs
ALM1	5	10	5	11k	66k	166k	14	86	1
ALM2	6	10	5	111k	666k	1.6M	22	387	5
ALM3	6	10	10	111k	1.2M	3.3M	29	1638	5
ALM4	5	24	5	346k	2.1M	5.2M	33	856	8
ALM5	4	64	12	266k	3.4M	9.6M	18	1195	8
ALM6	4	120	5	1.7M	10.4M	26.1M	18	1470	16
ALM7	4	120	10	1.7M	19.1M	52.2M	19	8465	16

# Warm-start strategies

A **warm-start strategy** uses the solution to a problem instance to initialise the next problem.

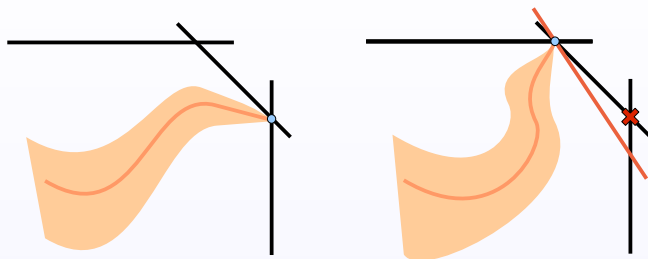
- ▶ Important if we are solving a sequence of problems
- ▶ Often we may expect that the solution to one problem is close to the solution of the next
- ▶ An advanced starting point may lead to reduced computational time than solving the problem from scratch

## Warm-start with IPMs



The solution of a problem is arbitrarily close to a [vertex](#):

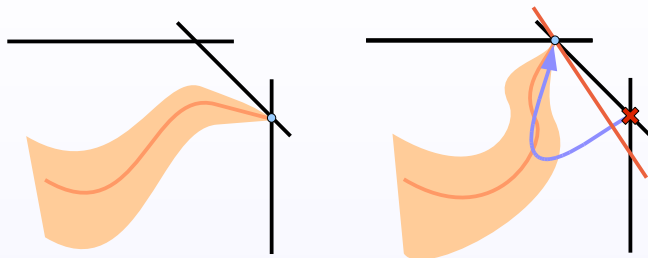
## Warm-start with IPMs



The solution of a problem is arbitrarily close to a **vertex**:

- ▶ **worst possible** starting point

## Warm-start with IPMs



The solution of a problem is arbitrarily close to a **vertex**:

- ▶ **worst possible** starting point
- ▶ some iterations to approach the new central path
- ▶ some iterations for optimality

# Warm-start for stochastic problems

## Observation:

Very detailed event trees provide a fine-grained solution to a problem that could have been solved more coarsely with a much smaller tree.

## Reduced event tree:

Use the solution to a smaller instance of the problem to generate a warm-start point.

## Main assumptions:

- ▶ No knowledge on the underlying stochastic processes
- ▶ Required to solve an instance with a specific tree
- ▶ We generate and solve the deterministic equivalent

## Scenario distance and representative scenarios

Distance between two nodes at period  $t$ :

$$d(n^{i_t}, n^{j_t}) = \|T^{i_t} - T^{j_t}\| + \|W^{i_t} - W^{j_t}\| + \|h^{i_t} - h^{j_t}\| + \|q^{i_t} - q^{j_t}\|$$

Distance between two scenarios:

$$D(s_i, s_j) = \sum_{t=1}^T d(n^{i_t}, n^{j_t}), \quad i_t \in s_i, j_t \in s_j$$

# Scenario distance and representative scenarios

Distance between two nodes at period  $t$ :

$$d(n^{it}, n^{jt}) = \|T^{it} - T^{jt}\| + \|W^{it} - W^{jt}\| + \|h^{it} - h^{jt}\| + \|q^{it} - q^{jt}\|$$

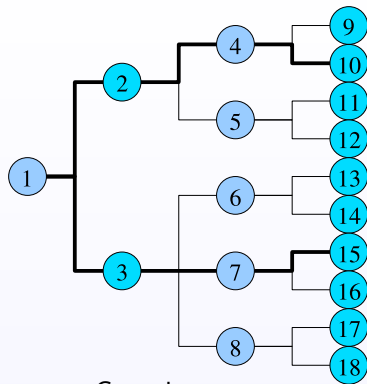
Distance between two scenarios:

$$D(s_i, s_j) = \sum_{t=1}^T d(n^{it}, n^{jt}), \quad i_t \in s_i, j_t \in s_j$$

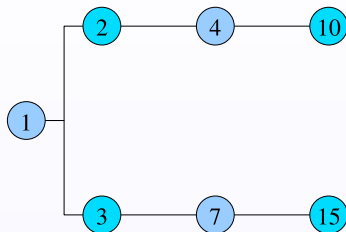
**Representative scenario  $s^*$**  is the one that minimizes the weighted distance from an average scenario  $\bar{s}$ :

$$s^* = s_k, \quad k = \arg \min_{i \in S} (1 - p_i) D(s_i, \bar{s})$$

## Reduced-tree warm-start technique



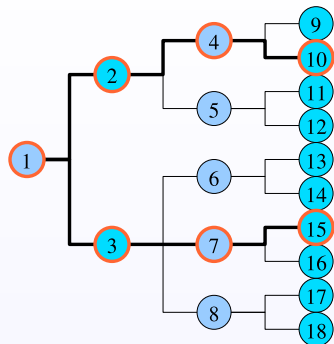
Complete tree



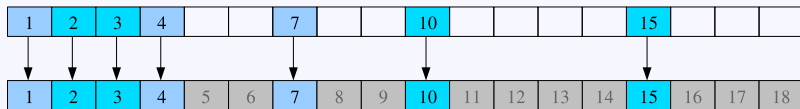
Reduced tree

1. Solve the problem with a reduced scenario tree
2. Expand the solution found to construct a starting point for the complete formulation
3. Solve the problem with the complete scenario tree

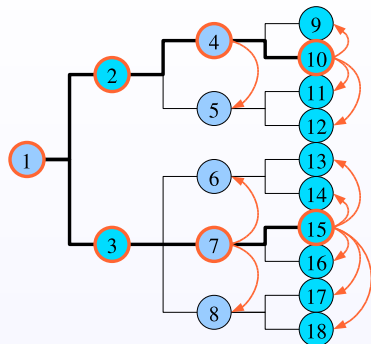
# Construction of the warm-start iterate



Nodes in the reduced tree:  
the solution is already available

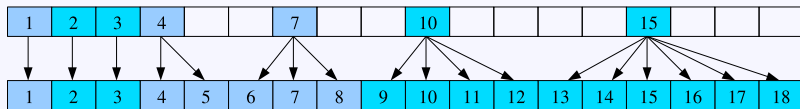


# Construction of the warm-start iterate



Nodes in the reduced tree:  
the solution is already available

Remaining nodes:  
copy the solution from the  
corresponding reduced-tree node



## Numerical results with HOPDM

Problem data			Cold start		Warm start	
Name	Stgs	Scens	Iters	Seconds	Iters	Seconds
fxm2-16	2	16	22	1.2	13	1.0
fxm3-6	3	36	30	1.5	17	1.3
fxm3-16	3	256	40	31.1	20	20.7
fxm4-6	4	216	30	8.2	22	8.3
fxm4-16	4	4096	41	218.3	27	182.6
pltexpA3-16	3	256	26	153.8	14	87.8
pltexpA4-6	4	216	36	55.8	16	27.5
pltexpA5-6	5	1296	81	772.0	30	311.5
storm27	2	27	41	95.4	22	53.2
storm125	2	125	73	107.3	36	69.1
storm1000	2	1000	107	1498.3	45	831.5

## Numerical results with OOPS (4 processors)

Problem data			Cold start		Warm start	
Name	Stgs	Scens	Iters	Seconds	Iters	Seconds
mnx-200	2	200	13	4.6	7	3.5
mnx-800	2	800	17	18.8	10	10.7
mnx-1600	2	1600	19	50.3	10	31.4
jlg-200	2	200	45	49.9	17	20.7
jlg-800	2	800	29	130.5	10	50.1
jlg-1600	2	1600	35	286.1	14	129.7
mgntA-100	2	100	28	76.9	14	51.6
mgntA-200	2	200	50	256.4	34	195.3
mgntA-400	2	400	40	410.9	14	181.6
mgntB-100	2	100	23	137.5	14	103.9
mgntB-200	2	200	25	284.2	8	140.5
mgntB-400	2	400	29	605.5	7	211.6

# Conclusions

- ▶ Structure can be exploited both at the linear algebra level and algorithmically
- ▶ OOPS provides an efficient implementation of a structure-exploiting parallel software
- ▶ Reduced tree solutions contain valuable information to construct a good warm-start iterate
- ▶ IPMs can be used successfully warm-started

# References

Gondzio, Grothey, Solving non-linear portfolio optimization problems with the primal–dual interior point method. EJOR 2006.

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