ADMM-based Unit and Time Decomposition for Price Arbitrage by Cooperative Price-Maker Electricity Storage Units

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Joint work with Miguel F. Anjos¹ and James R. Cruise²

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19/05/2022

Introduction - The Electric Power System

• Electric power systems must be in close balance at any time:

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- Traditional approach: Keep flexible power plants on call:
 - 1. Hydro
 - 2. Gas
 - 3. Diesel, coal and biomass
- Challenges:
 - 1. Hydro resources are limited.
 - 2. Other resources produce carbon emissions.

Introduction - Increasing Flexibility Needs

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Image: A matrix and a matrix

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This increases flexibility needs for two reasons:

- 1. Higher shares of wind, solar and/or nuclear.
- 2. Reduction of the sources that currently provide a substantial amount of flexibility: gas, diesel and coal.

1. A well-connected grid

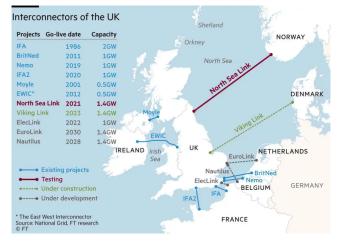


Figure: Interconnectors between the UK and other European countries: operational (blue and purple) and under construction or planned (other).

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2. Electric energy storage



Figure: A pumped-storage hydro station.



Figure: A lithium-ion battery.

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3. Demand-side response





Figure: The canal network and an electric vehicle.





Figure: A fridge and an electric water heater.

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4. Fossil fuels + carbon capture and storage

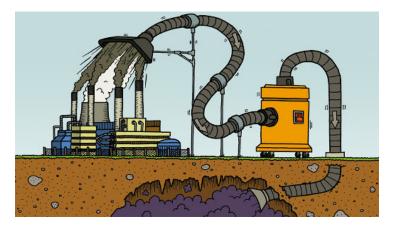


Figure: Liberal representation of carbon capture and storage.

The future flexibility requirements will probably be met from many sources. We study one of them:

Electric energy storage

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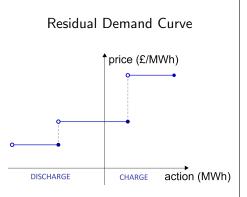
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Research Question: How should we control multiple price-maker electric energy storage units that cooperate for price arbitrage?

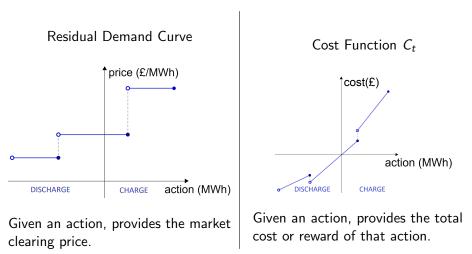
We extend previous work on a single unit by Cruise et al. (2019), and our previous work on two units (Anjos, Cruise, SV (2020)).

Modelling Price-Makers



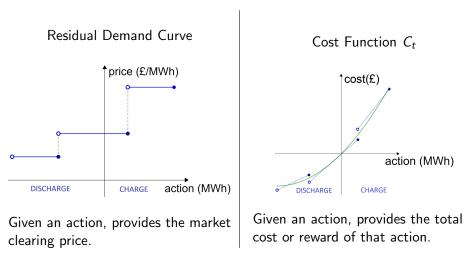
Given an action, provides the market clearing price.

Modelling Price-Makers



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Modelling Price-Makers



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Key assumption: C_t is convex.
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Let $x_{j,t} \in \mathbb{R}$ be the action taken by unit $j \in S$ at time $t \in \mathcal{T}$.

 $\underset{x_{j,t}}{\mathsf{Minimize}}$

 $\sum_{t\in\mathcal{T}} C_t \Big(\sum_{j\in\mathcal{S}} x_{j,t}\Big)$

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We assume:

- Discrete Time
- Deterministic Prices

A. Solà Vilalta (UoE)

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*The model can account for round-trip efficiencies, leakage and negative electricity prices.

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 $\underset{x_{j,t}}{\mathsf{Minimize}} \qquad \sum_{t \in \mathcal{T}} \mathbf{0}$

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Minimize $x_{j,t}, z_t$

$$\sum_{t\in\mathcal{T}}C_t(\mathbf{z}_t)$$

subject to

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$$\begin{array}{ll} \underset{x_{j,t},z_{t}}{\text{Minimize}} & \sum_{t \in \mathcal{T}} \left[\mathcal{C}_{t}(z_{t}) + \nu_{t} \Big(\sum_{j \in \mathcal{S}} x_{j,t} - z_{t} \Big) + \frac{\gamma}{2} \Big(\sum_{j \in \mathcal{S}} x_{j,t} - z_{t} \Big)^{2} \right] \\ \text{subject to} & -P_{j} \leq x_{j,t} \leq P_{j} & \forall j \in \mathcal{S}, \ \forall t \in \mathcal{T} \\ & 0 \leq \bar{S}_{j,0} + \sum_{l=1}^{t} x_{j,l} \leq E_{j} & \forall j \in \mathcal{S}, \ \forall t \in \mathcal{T} \\ & \bar{S}_{j,0} + \sum_{l=1}^{T} x_{j,l} = \bar{S}_{j,\mathcal{T}} & \forall j \in \mathcal{S} \end{array}$$

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Convexity and linear constraints imply strong duality $$\downarrow$$ Use the Alternating Direction Method of Multipliers (ADMM)

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This results in N + 1 1-Unit Subproblems:

- A subproblem for every storage unit.
- A subproblem for the auxiliary variable z.

They are solved iteratively until the linking constraints are satisfied.

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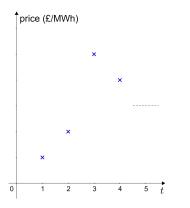
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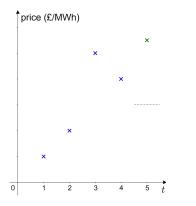
Ex: Consider a price-taker storage unit with energy capacity E = 2 and power rate P = 1 with initial SoC $\overline{S}_0 = 0$. Assume prices are given by:



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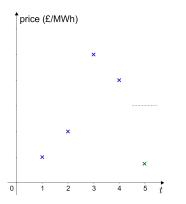
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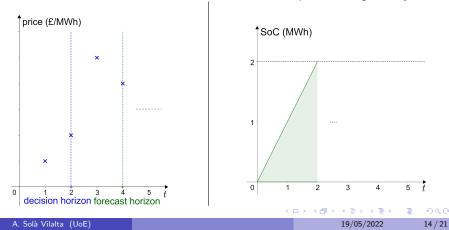
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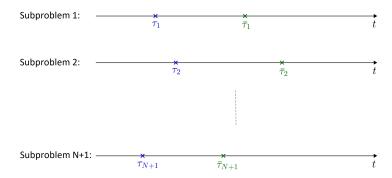
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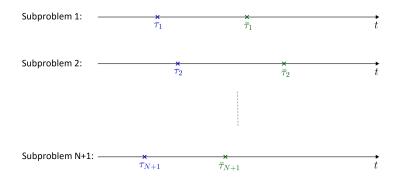
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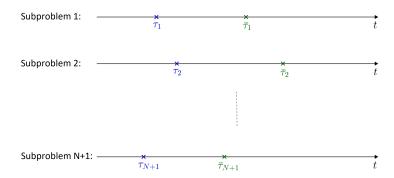
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Image: A matrix and a matrix

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A: Pick the longest forecast horizon and solve all subproblems until then. \downarrow It is the forecast horizon of the unit with the largest energy-to-power ratio.

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Results

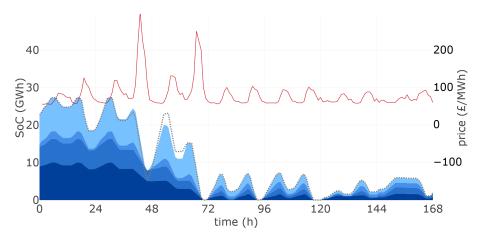


Figure: Left axis: Stacked SoC of unit 1 [very dark blue], unit 2 [dark blue], unit 3 [light blue] and unit 4 [very light blue]. SoC of aggregated unit [dotted grey line]. Right axis: Electricity market-clearing prices [red].

Computational Performance I

Instance	Anjos et al. (2020)	Anjos et al. (2021)	Gap
Oct '19	$-6.231 imes10^{6}$	$-6.230 imes10^{6}$	0.02%
Nov '19	$-5.689 imes10^{6}$	$-5.688 imes10^{6}$	0.02%
Dec '19		$-7.091 imes10^{6}$	
Jan '20	$-5.988 imes10^{6}$	$-5.986 imes10^6$	0.03%
Feb '20	$-6.558 imes10^{6}$	$-6.953 imes10^{6}$	-6.03%
Mar '20	$6.384 imes10^{6}$	$6.378 imes10^{6}$	0.09%

Table: Objective function values and gap between the methods in Anjos, Cruise, SV (2020) and Anjos, Cruise, SV (2021). One month instances of 2-Unit Problems.

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Computational Performance II

Month	Anjos et al. (2020)	Anjos et al. (2021)
Oct '19	91s	0.83s
Nov '19	89s	0.54s
Dec '19		0.45s
Jan '20	112s	0.56s
Feb '20	83s	0.28s
Mar '20	103s	0.19s

Table: Computational time comparison of the solution methods in Anjos, Cruise, SV (2020) and Anjos, Cruise, SV (2021). One month instances of 2-Unit Problems.

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Computational Performance III

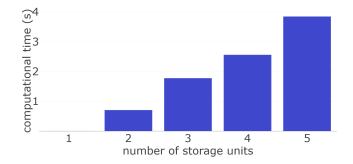


Figure: Average computational time of one month instances for different number of storage units.

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- Two orders of magnitude computational time reduction compared to the algorithm in Anjos, Cruise, SV (2020) in 2-unit instances, with minor solution quality losses (< 0.1%).
- Linear scaling of computational time w.r.t. the number of units.
- The energy-to-power ratio of storage units plays a crucial role.

References

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- Anjos MF, Cruise JR, Solà Vilalta A (2021) ADMM-based Unit and Time Decomposition for Price Arbitrage by Cooperative Price-Maker Electricity Storage Units. Working paper.
- Cruise JR, Flatley L, Gibbens RJ, Zachary S (2019) Control of energy storage with market impact: Lagrangian approach and horizons. Operations Research 67(1):1–9.

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